

# Object-Oriented Design Case Studies with Patterns & C++

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## Case Studies Using Patterns

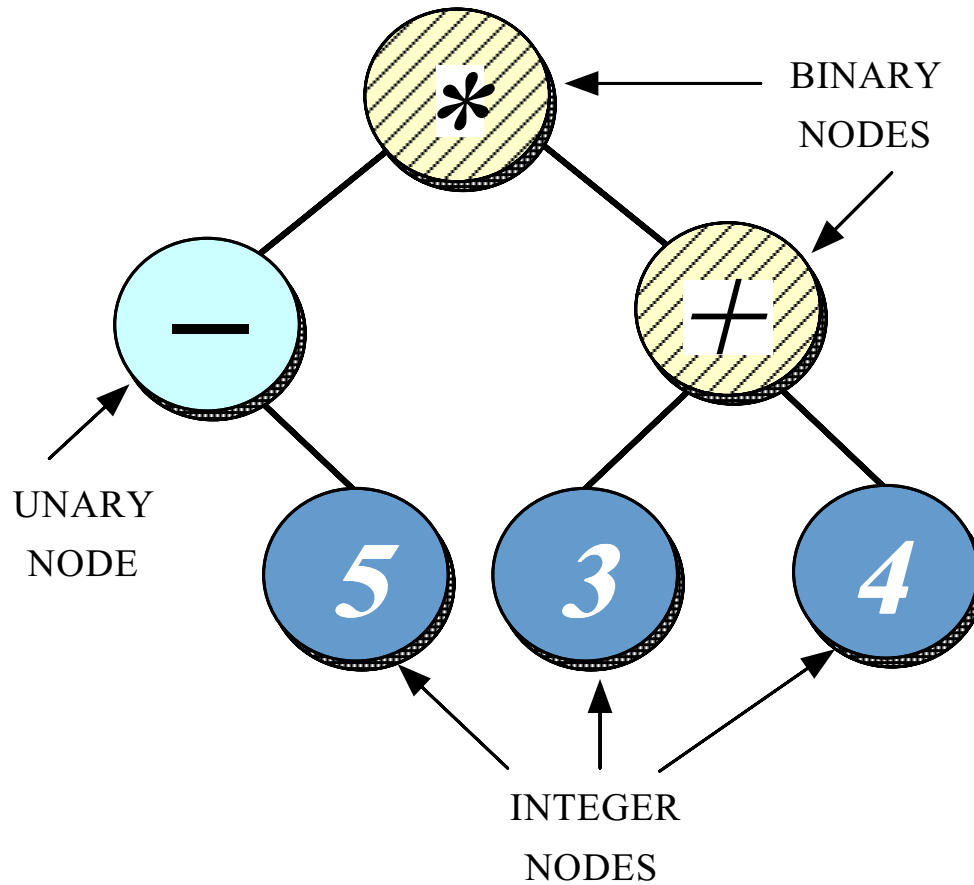
- The following slides describe several case studies using C++ & patterns to build highly extensible software
- The examples include
  1. Expression Tree
    - e.g., Adapter, Factory, Bridge
  2. System Sort
    - e.g., Facade, Adapter, Iterator, Singleton, Factory Method, Strategy, Bridge
  3. Sort Verifier
    - e.g., Strategy, Factory Method, Facade, Iterator, Singleton

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## Case Study: Expression Tree Evaluator

- The following inheritance & dynamic binding example constructs *expression trees*
  - Expression trees consist of nodes containing operators & operands
    - \* Operators have different *precedence levels*, different *associativities*, & different *arities*, e.g.,
      - Multiplication takes precedence over addition
      - The multiplication operator has two arguments, whereas unary minus operator has only one
    - \* Operands are integers, doubles, variables, etc.
      - We'll just handle integers in this example . . .

# Expression Tree Diagram



## Expression Tree Behavior

- *Expression trees*
  - Trees may be “evaluated” via different traversals
    - \* e.g., in-order, post-order, pre-order, level-order
  - The evaluation step may perform various operations, e.g.,
    - \* Traverse & print the expression tree
    - \* Return the “value” of the expression tree
    - \* Generate code
    - \* Perform semantic analysis

## Algorithmic Version

- A typical algorithmic method for implementing expression trees involves using a struct/union to represent data structure, e.g.,

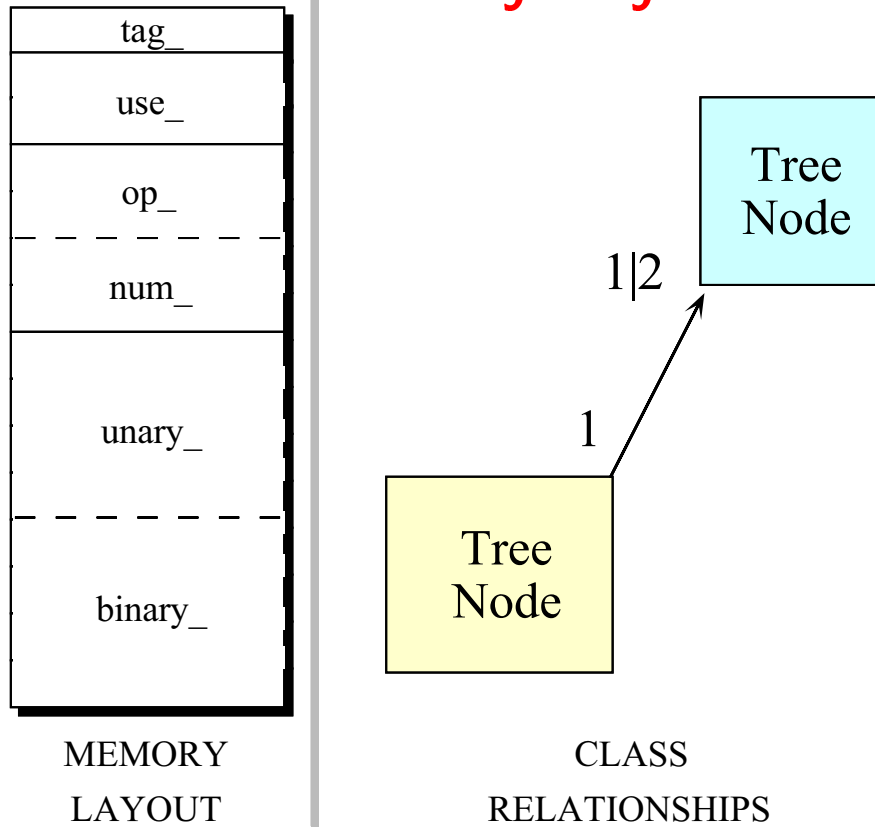
```

typedef struct Tree_Node Tree_Node;
struct Tree_Node {
    enum { NUM, UNARY, BINARY } tag;
    short use; /* reference count */
    union {
        char op[2];
        int num;
    } o;
#define num_o.o.num
#define op_o.o.op
    union {
        Tree_Node *unary;
        struct { Tree_Node *l, *r; } binary;
    } c;
#define unary_c.unary
#define binary_c.binary
};

```



## Memory Layout of Algorithmic Version



- Here's the memory layout of a struct Tree\_Node object

## Print\_Tree Function

- A typical algorithmic implementation use a switch statement & a recursive function to build & evaluate a tree, e.g.,

```

void print_tree (TreeNode *root) {
    switch (root->tag-) {
        case NUM: printf ("%d", root->num-);
            break;
        case UNARY:
            printf ("%s", root->op-[0]);
            print_tree (root->unary-);
            break;
        case BINARY:
            printf ("%s", root->op-[0]);
            print_tree (root->binary-.l-);
            print_tree (root->binary-.r-);
            break;
        default:
            printf ("error, unknown type\n");
    }
}

```





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## Limitations with Algorithmic Approach

- Problems or limitations with the typical algorithmic approach include
  - Little or no use of encapsulation
  - Incomplete modeling of the application domain, which results in
    1. Tight coupling between nodes & edges in union representation
    2. Complexity being in *algorithms* rather than the *data structures*
      - e.g., switch statements are used to select between various types of nodes in the expression trees
      - Compare with binary search!
    3. Data structures are “passive” & functions do most processing work explicitly

## More Limitations with Algorithmic Approach

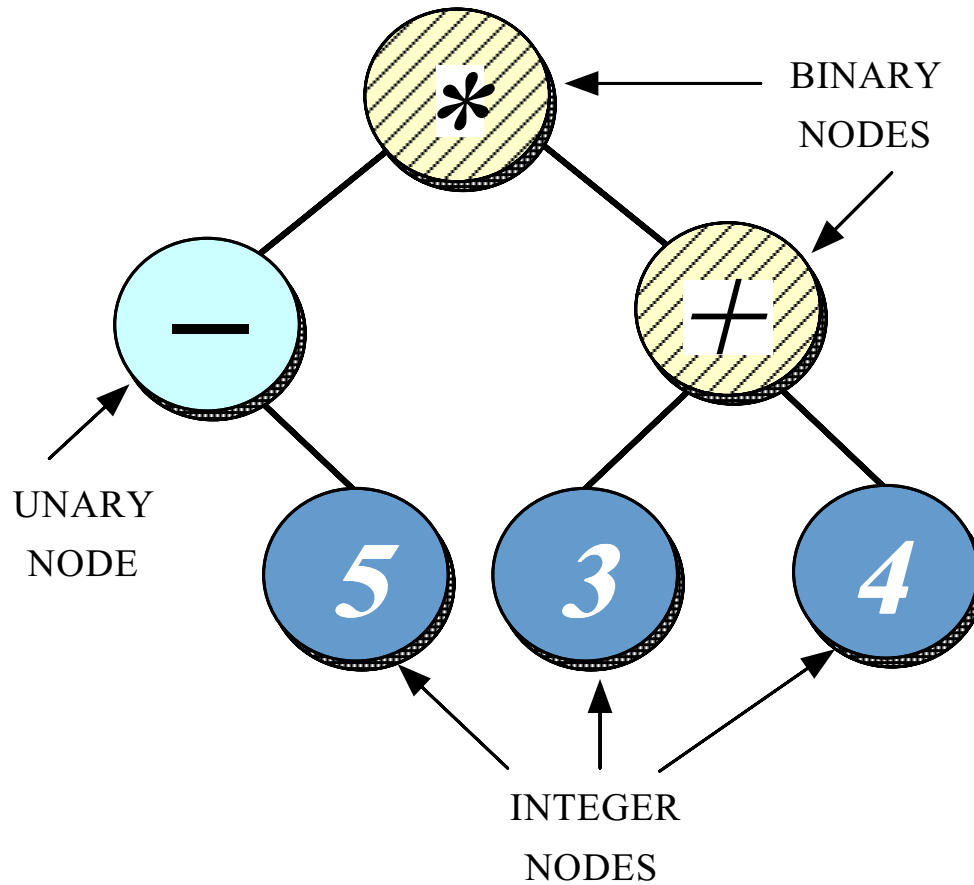
- The program organization makes it difficult to extend, e.g.,
  - Any small changes will ripple through the entire design & implementation
    - \* e.g., see the “ternary” extension below
  - Easy to make mistakes switching on type tags . . .
- Solution wastes space by making worst-case assumptions wrt structs & unions
  - This is not essential, but typically occurs
  - Note that this problem becomes worse the bigger the size of the largest item becomes!

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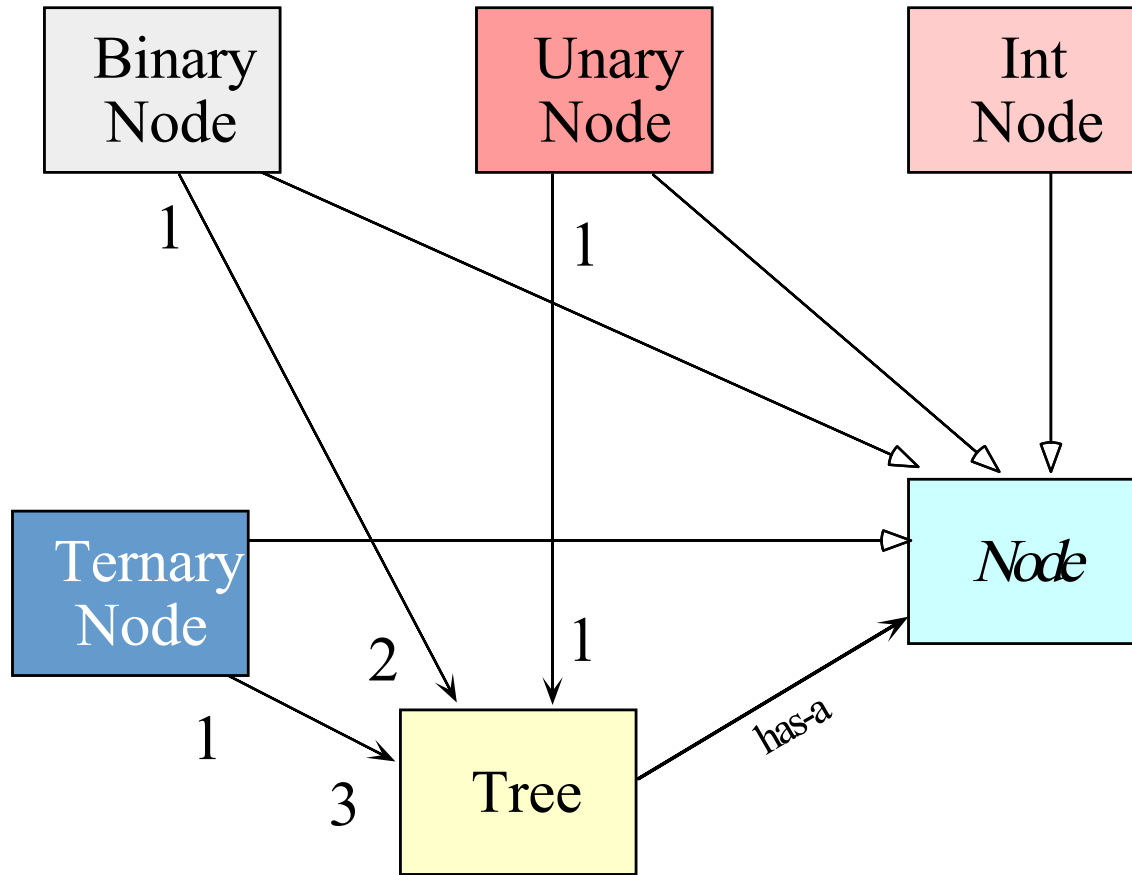
## OO Alternative

- Contrast previous algorithmic approach with an object-oriented decomposition for the same problem:
  - Start with OO modeling of the “expression tree” application domain, e.g., go back to original picture
  - Discover several classes involved:
    - \* class Node: base class that describes expression tree vertices:
      - class Int\_Node: used for implicitly converting int to Tree node
      - class Unary\_Node: handles unary operators, e.g., -10, +10, !a
      - class Binary\_Node: handles binary operators, e.g., a + b, 10 - 30
    - \* class Tree: “glue” code that describes expression-tree edges, *i.e.*, relations between Nodes
  - Note, these classes model entities in the application domain
    - \* *i.e.*, nodes & edges (vertices & arcs)

# Expression Tree Diagram



## Relationships Between Tree & Node Classes



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## Design Patterns in the Expression Tree Program

- Factory
  - Centralize the assembly of resources necessary to create an object
    - \* e.g., decouple Node subclass initialization from use
- Bridge
  - Decouple an abstraction from its implementation so that the two can vary independently
    - \* e.g., printing contents of a subtree and managing memory
- Adapter
  - Convert the interface of a class into another interface clients expect
    - \* e.g., make Tree conform C++ iostreams

## C++ Node Interface

```
class Tree; // Forward declaration

// Describes the Tree vertices
class Node {
friend class Tree;
protected: // Only visible to derived classes
    Node (): use_ (1) {}

    /* pure */ virtual void print (std::ostream &) const = 0;

    // Important to make destructor virtual!
    virtual ~Node ();
private:
    int use_; // Reference counter.
};
```

## C++ Tree Interface

```
#include "Node.h"
// Bridge class that describes the Tree edges and
// acts as a Factory.
class Tree {
public:
    // Factory operations
    Tree (int);
    Tree (const string &, Tree &);
    Tree (const string &, Tree &, Tree &);
    Tree (const Tree &);
    void operator= (const Tree &);
    ~Tree ();
    void print (std::ostream &) const;
private:
    Node *node_; // pointer to a rooted subtree
```





## C++ Int\_Node Interface

```
#include "Node.h"

class Int_Node : public Node {
public:
    Int_Node (int k);
    virtual void print (std::ostream &stream) const;
private:
    int num_; // operand value.
};
```

## C++ Unary\_Node Interface

```
#include "Node.h"

class Unary_Node : public Node {
public:
    Unary_Node (const string &op, const Tree &t);
    virtual void print (std::ostream &stream) const;
private:
    string operation_;
    Tree operand_;
};
```

## C++ Binary\_Node Interface

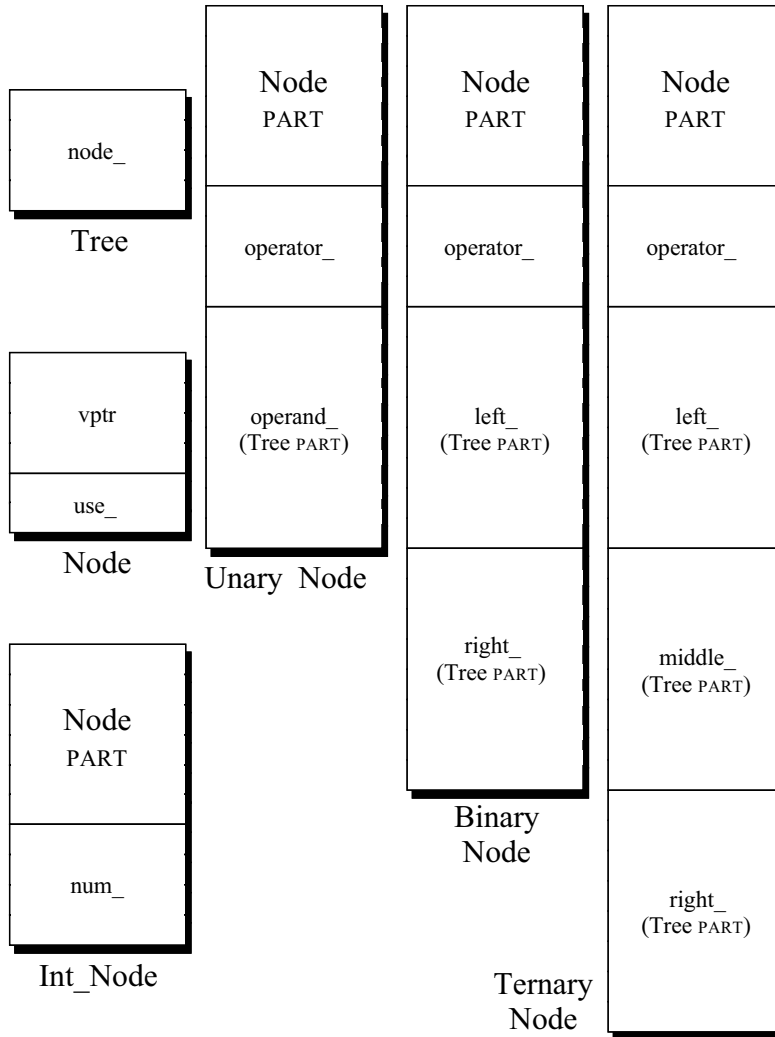
```
#include "Node.h"

class Binary_Node : public Node {
public:
    Binary_Node (const string &op,
                const Tree &t1,
                const Tree &t2);

    virtual void print (std::ostream &s) const;

private:
    const string operation_;
    Tree left_;
    Tree right_;
};
```

# Memory Layout for C++ Version



- Memory layouts for different subclasses of Node

## C++ Int\_Node Implementations

```
#include "Int_Node.h"

Int_Node::Int_Node (int k): num_ (k) { }

void Int_Node::print (std::ostream &stream) const {
    stream << this->num_;
}
```

## C++ Unary\_Node Implementations

```
#include "Unary_Node.h"

Unary_Node::Unary_Node (const string &op, const Tree &t1)
    : operation_ (op), operand_ (t1) { }

void Unary_Node::print (std::ostream &stream) const {
    stream << "(" << this->operation_ <<
        << this->operand_ // recursive call!
        << ")";
}
```

## C++ Binary\_Node Implementation

```
#include "Binary_Node.h"

Binary_Node::Binary_Node (const string &op,
                          const Tree &t1,
                          const Tree &t2):
    operation_ (op), left_ (t1), right_ (t2) {}

void Binary_Node::print (std::ostream &stream) const {
    stream << "(" << this->left_ // recursive call
          << " " << this->operation_
          << " " << this->right_ // recursive call
          << ")";
}
```

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## Initializing the Node Subclasses

- *Problem*
  - How to ensure the Node subclasses are initialized properly
- *Forces*
  - There are different types of Node subclasses
    - \* e.g., take different number & type of arguments
  - We want to centralize initialization in one place because it is likely to change . . .
- *Solution*
  - Use a *Factory* pattern to initialize the Node subclasses

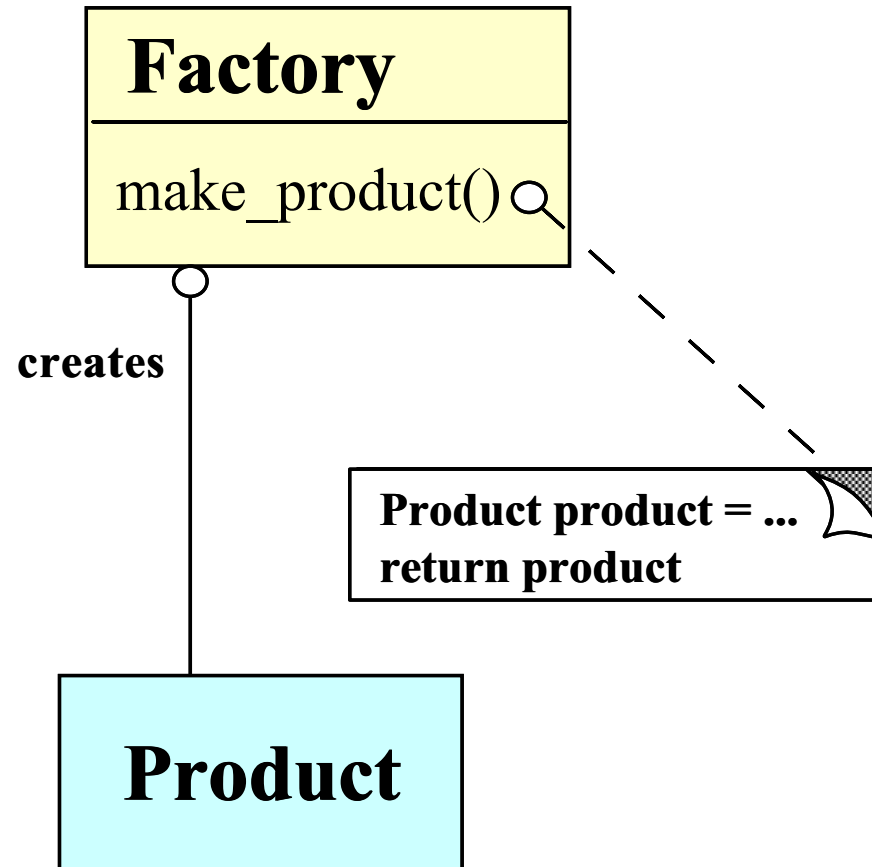


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## The Factory Pattern

- *Intent*
  - *Centralize the assembly of resources necessary to create an object*
    - \* Decouple object creation from object use by localizing creation knowledge
- This pattern resolves the following forces:
  - Decouple initialization of the Node subclasses from their subsequent use
  - Makes it easier to change or add new Node subclasses later on
    - \* e.g., Ternary nodes . . .
- A generalization of the GoF Factory Method pattern

## Structure of the Factory Pattern



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## Using the Factory Pattern

- The Factory pattern is used by the Tree class to initialize Node subclasses:

```
Tree::Tree (int num)
: node_ (new Int_Node (num)) {}
```

```
Tree::Tree (const string &op, const Tree &t)
: node_ (new Unary_Node (op, t)) {}
```

```
Tree::Tree (const string &op,
            const Tree &t1,
            const Tree &t2)
: node_ (new Binary_Node (op, t1, t2)) {}
```

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## Printing Subtrees

- *Problem*
  - How do we print subtrees without revealing their types?
- *Forces*
  - The `Node` subclass should be hidden within the `Tree` instances
  - We don't want to become dependent on the use of `Nodes`, inheritance, & dynamic binding, etc.
  - We don't want to expose dynamic memory management details to application developers
- *Solution*
  - Use the *Bridge* pattern to shield the use of inheritance & dynamic binding

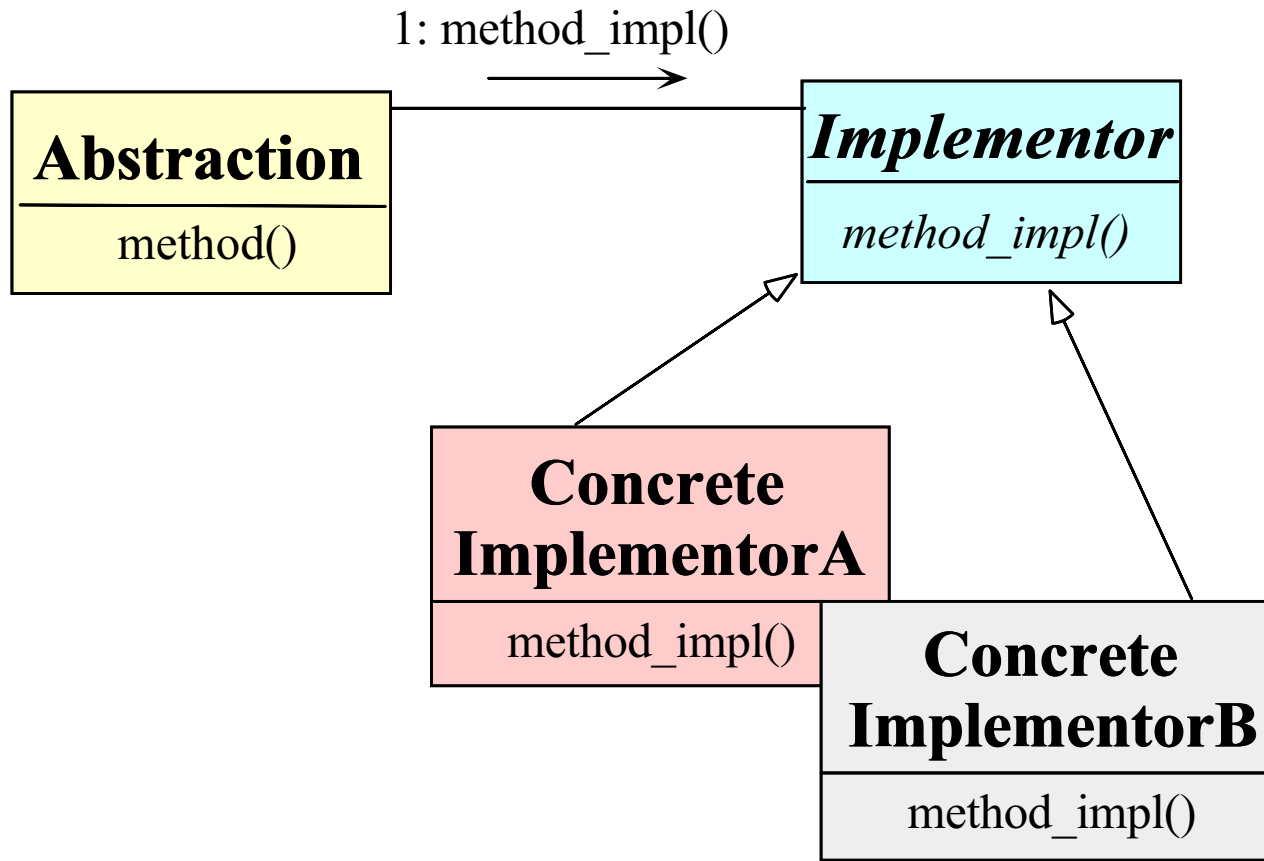


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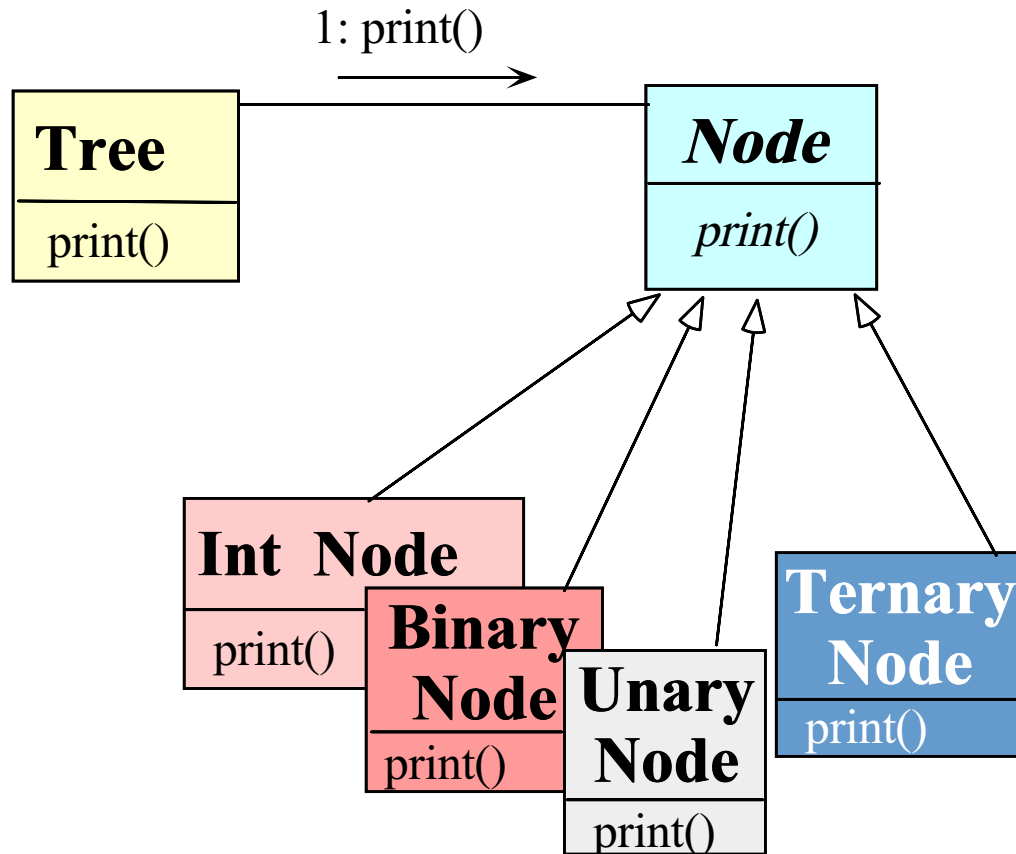
## The Bridge Pattern

- *Intent*
  - *Decouple an abstraction from its implementation so that the two can vary independently*
- This pattern resolves the following forces that arise when building extensible software with C++
  1. *How to provide a stable, uniform interface that is both closed & open, i.e.,*
    - interface is closed to prevent direct code changes
    - Implementation is open to allow extensibility
  2. *How to manage dynamic memory more transparently & robustly*
  3. *How to simplify the implementation of operator<<*

# Structure of the Bridge Pattern



# Using the Bridge Pattern



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## Illustrating the Bridge Pattern in C++

- The Bridge pattern is used for printing expression trees:

```
void Tree::print (std::ostream &os) const {  
    this->node_>print (os);  
}
```

- Note how this pattern decouples the Tree interface for printing from the Node subclass implementation
  - *i.e.*, the Tree interface is *fixed*, whereas the Node implementation varies
  - However, clients need not be concerned about the variation . . .



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## Integrating with C++ I/O Streams

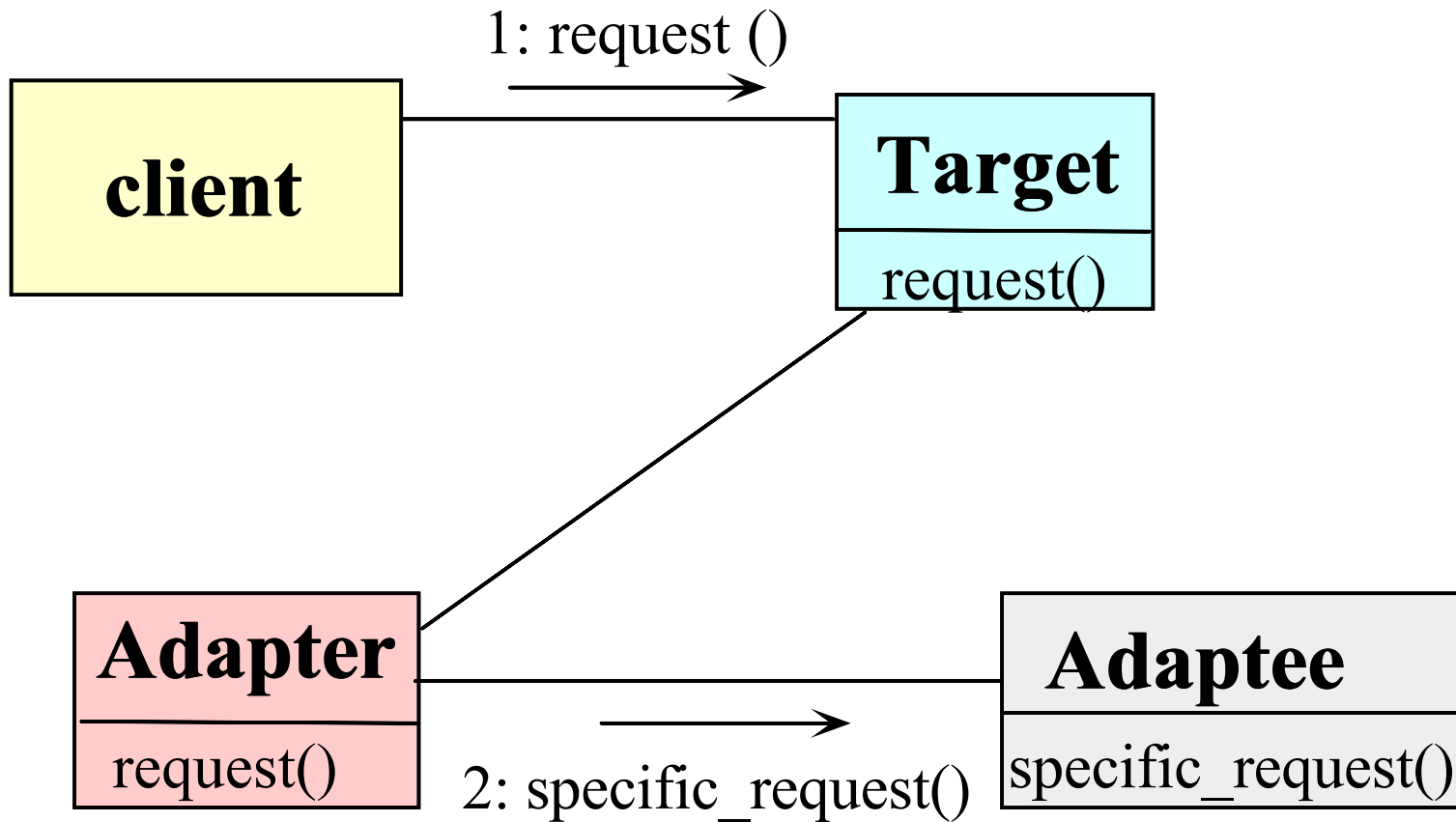
- *Problem*
  - Our Tree interface uses a print method, but most C++ programmers expect to use I/O Streams
- *Forces*
  - Want to integrate our existing C++ Tree class into the I/O Stream paradigm without modifying our class or C++ I/O
- *Solution*
  - Use the *Adapter* pattern to integrate Tree with I/O Streams

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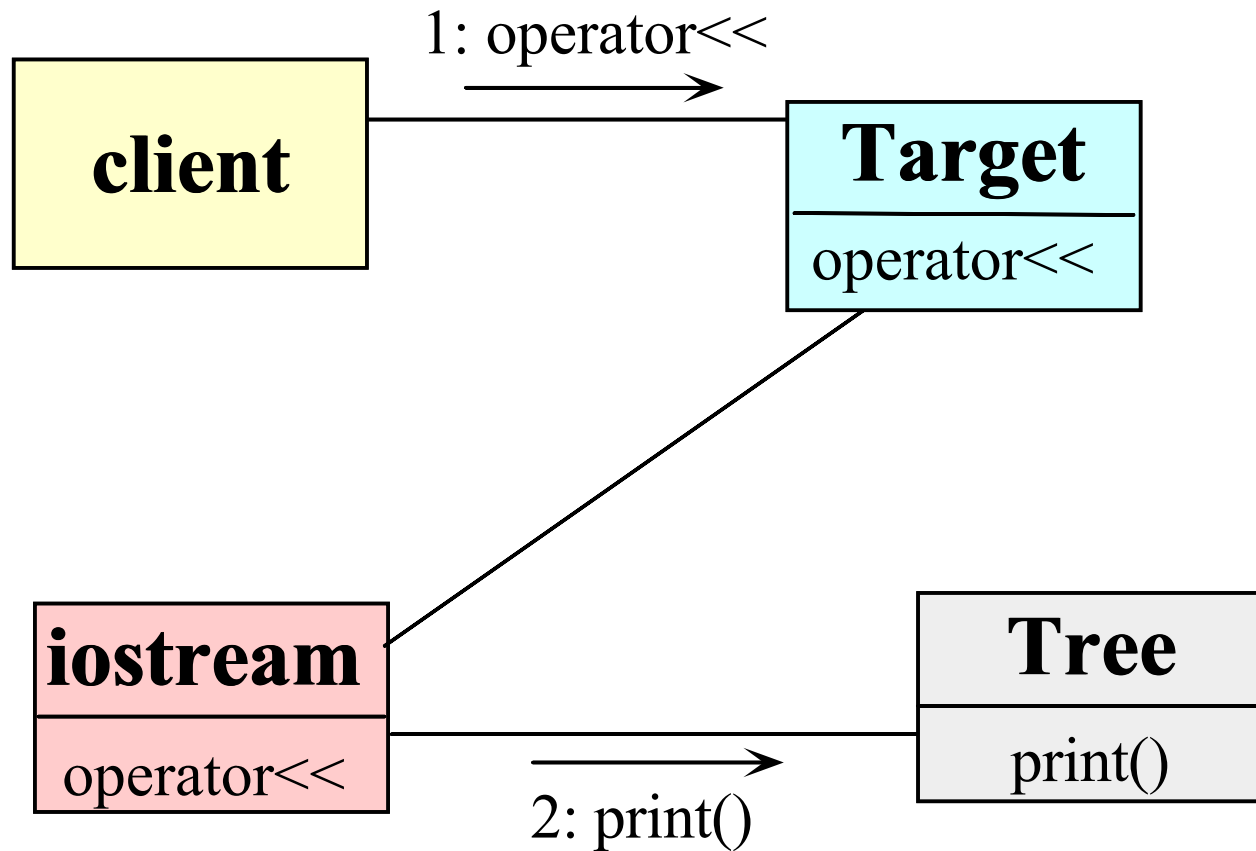
## The Adapter Pattern

- *Intent*
  - Convert the interface of a class into another interface client expects
    - \* Adapter lets classes work together that couldn't otherwise because of incompatible interfaces
- This pattern resolves the following force:
  1. How to transparently integrate the `Tree` with the `C++ istd::ostream` operators

## Structure of the Adapter Pattern



# Using the Adapter Pattern



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## Using the Adapter Pattern

- The Adapter pattern is used to integrate with C++ I/O Streams

```
std::ostream &operator<< (std::ostream &s, const Tree &ttree) {
    ttree.print (s);
    // This triggers Node * virtual call via
    // ttree.node_>print (s), which is
    // implemented as the following:
    // (*tree.node_>vptr[1]) (tree.node_, s);
    return s;
}
```

- Note how the C++ code shown above uses I/O streams to “adapt” the Tree interface . . .

## C++ Tree Implementation

- Reference counting via the “counted body” idiom

```
Tree::Tree (const Tree &t): node_ (t.node_) {
    ++this->node_->use_; // Sharing, ref-counting.
}

void Tree::operator= (const Tree &t) {
    if (this == &t) return;
    // order important here!
    ++t.node_->use_;
    --this->node_->use_;
    if (this->node_->use_ == 0)
        delete this->node_;
    this->node_ = t.node_;
}
```

## C++ Tree Implementation (cont'd)

```
Tree::~Tree () {  
    // Ref-counting, garbage collection  
    --this->node_->use_;  
    if (this->node_->use_<= 0)  
        delete this->node_;  
}
```

## C++ Main Program

```
#include <istream>
#include <ostream>
#include "Tree.h"

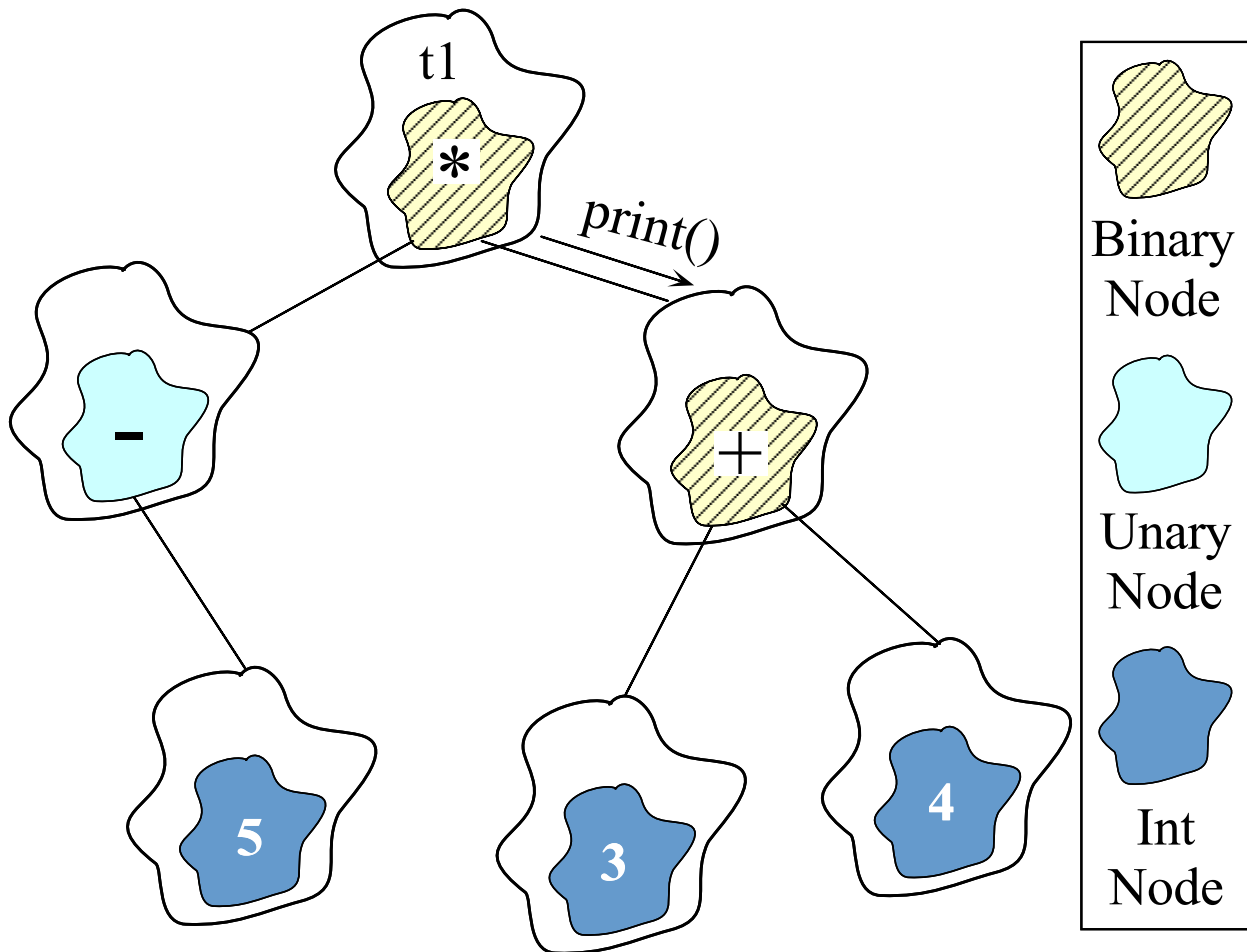
int main (int, char *[]) {
    const Tree t1 = Tree ("*", Tree ("-", 5),
                          Tree ("+", 3, 4));
    cout << t1 << endl; // prints ((-5) * (3 + 4))
    const Tree t2 = Tree ("*", t1, t1);

    // prints (((-5) * (3 + 4)) * ((-5) * (3 + 4))).
    cout << t2 << endl;

    return 0;
    // Destructors of t1 & t2 recursively
    // delete entire tree when leaving scope.
}
```

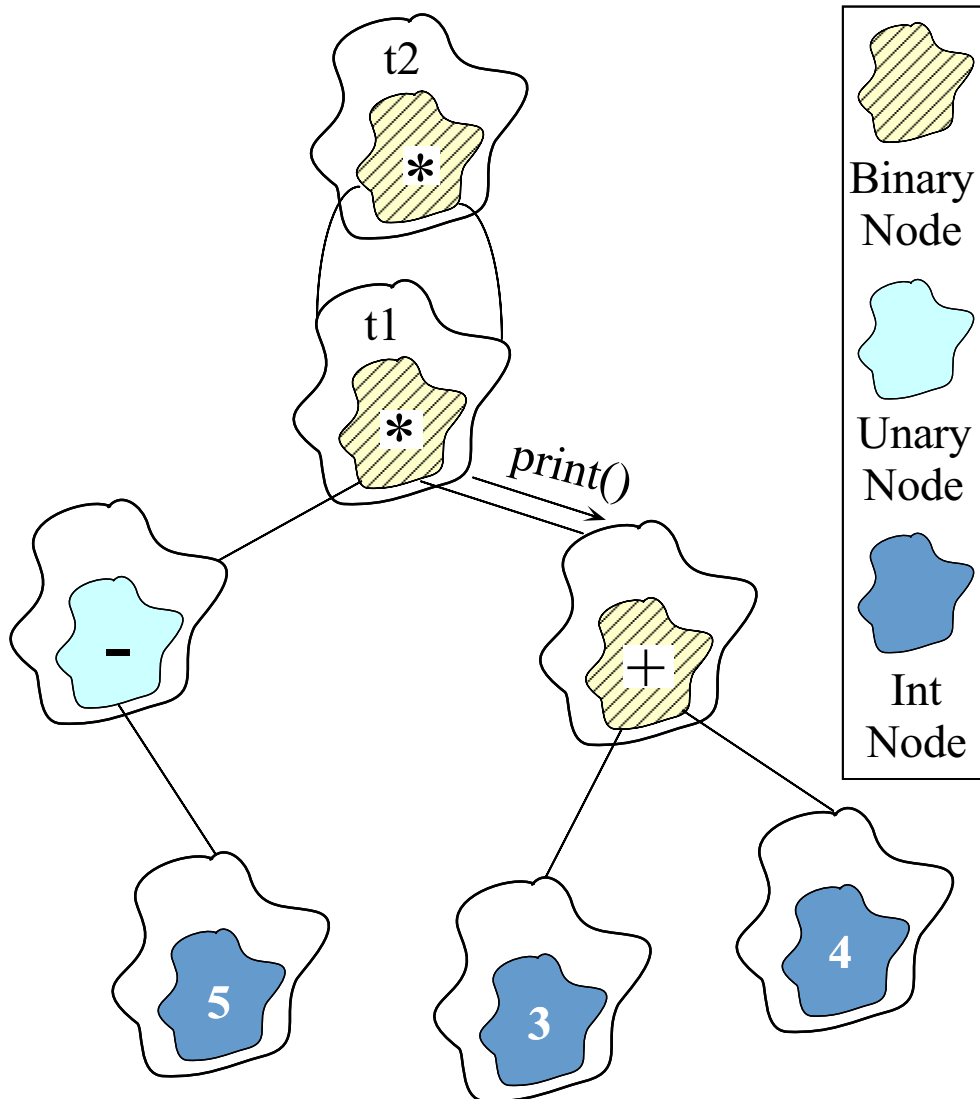


## Expression Tree Diagram 1



- Expression tree for `t1 = ((-5) * (3 + 4))`

## Expression Tree Diagram 2



- Expression tree for  $t2 = (t1 * t1)$

## Adding Ternary\_Nodes

- Extending the existing program to support ternary nodes is straightforward
  - *i.e.*, just derive new class Ternary\_Node to handle ternary operators, *e.g.*, `a == b ? c : d`, *etc.*

```
#include "Node.h"
class Ternary_Node : public Node {
public:
    Ternary_Node (const string &, const Tree &,
                 const Tree &, const Tree &);
    virtual void print (std::ostream &) const;
private:
    const string operation_;
    Tree left_, middle_, right_; };
```

## C++ Ternary\_Node Implementation

```
#include "Ternary_Node.h"
Ternary_Node::Ternary_Node (const string &op,
                             const Tree &a,
                             const Tree &b,
                             const Tree &c)
    : operation_ (op), left_ (a), middle_ (b),
      right_ (c) {}

void Ternary_Node::print (std::ostream &stream) const {
    stream << this->operation_ << "("
          << this->left_ // recursive call
          << ", " << this->middle_ // recursive call
          << ", " << this->right_ // recursive call
          << ")";
}
```

## C++ Ternary\_Node Implementation (cont'd)

```
// Modified class Tree Factory
class Tree {
// add 1 class constructor
public:
    Tree (const string &, const Tree &,
          const Tree &, const Tree &)
        : node_ (new Ternary_Node (op, l, m, r)) {}
// Same as before . . .
```

---

## Differences from Algorithmic Implementation

- On the other hand, modifying the original algorithmic approach requires changing (1) the original data structures, e.g.,

```
struct Tree_Node {
    enum {
        NUM, UNARY, BINARY, TERNARY
    } tag_; // same as before
    union {
        // same as before. But, add this:
        struct {
            Tree_Node *l_, *m_, *r_;
        } ternary_;
    } c;
#define ternary_ c.ternary_
};
```

## Differences from Algorithmic Implementation (cont'd)

- & (2) many parts of the code, e.g.,

```
void print_tree (Tree_Node *root) {  
    // same as before  
    case TERNARY: // must be TERNARY.  
        printf ("(");  
        print_tree (root->ternary_.l_);  
        printf ("%c", root->op_[0]);  
        print_tree (root->ternary_.m_);  
        printf ("%c", root->op_[1]);  
        print_tree (root->ternary_.r_);  
        printf (")"); break;  
    // same as before  
}
```

---

## Summary of Expression Tree Example

- OO version represents a more complete modeling of the application domain
  - e.g., splits data structures into modules that correspond to “objects” & relations in expression trees
- Use of C++ language features simplifies the design and facilitates extensibility
  - e.g., implementation follows directly from design
- Use of patterns helps to motivate, justify, & generalize design choices



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## Potential Problems with OO Design

- Solution is very “data structure rich”
  - e.g., requires configuration management to handle many headers & .cpp files!
- May be somewhat less efficient than original algorithmic approach
  - e.g., due to virtual function overhead
- In general, however, virtual functions may be no less inefficient than large switch statements or if/else chains . . .
- As a rule, be careful of micro vs. macro optimizations
  - *i.e.*, always profile your code!

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## Case Study: System Sort

- Develop a general-purpose system sort
  - It sorts lines of text from standard input and writes the result to standard output
  - e.g., the UNIX system sort
- In the following, we'll examine the primary forces that shape the design of this application
- For each force, we'll examine patterns that resolve it

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## External Behavior of System Sort

- A “line” is a sequence of characters terminated by a newline
- Default ordering is lexicographic by bytes in machine collating sequence (e.g., ASCII)
- The ordering is affected globally by the following options:
  - Ignore case (`-f`)
  - Sort numerically (`-n`)
  - Sort in reverse (`-r`)
  - Begin sorting at a specified field (`-k`)
  - Begin sorting at a specified column (`-c`)
- Your program need not sort files larger than main memory

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## High-level Forces

- Solution should be both time & space efficient
  - e.g., must use appropriate algorithms and data structures
  - Efficient I/O & memory management are particularly important
  - Our solution uses minimal dynamic binding (to avoid unnecessary overhead)
- Solution should leverage reusable components
  - e.g., `istd::ostreams`, `Array` & `Stack` classes, *etc.*
- Solution should yield reusable components
  - e.g., efficient input classes, generic sort routines, *etc.*

## Top-level Algorithmic View of the Solution

- Note the use of existing C++ mechanisms like I/O streams

```
// Reusable function:
// template <typename ARRAY> void sort (ARRAY &a);

int main (int argc, char *argv[])
{
    parse_args (argc, argv);
    Input input;

    cin >> input;
    sort (input);
    cout << input;
}
```

---

## Top-level Algorithmic View of the Solution (cont'd)

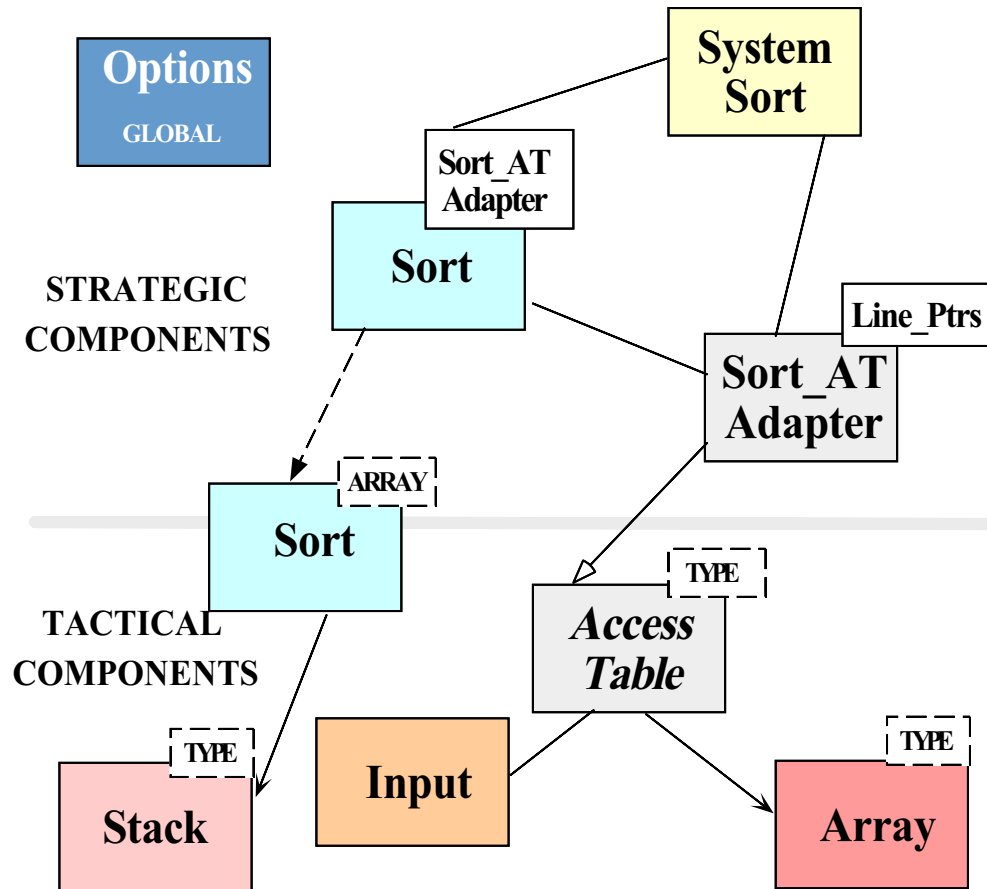
- Avoid the *grand mistake* of using top-level algorithmic view to structure the design . . .
  - Structure the design to resolve the forces!
  - Don't focus on algorithms or data, but instead look at the problem, its participants, & their interactions!

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## General OOD Solution Approach

- Identify the classes in the application/problem space & solution space
  - e.g., stack, array, input class, options, access table, sorts, etc.
- Recognize & apply common design patterns
  - e.g., Singleton, Factory, Adapter, Iterator
- Implement a framework to coordinate components
  - e.g., use C++ classes & parameterized types

# C++ Class Model





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## C++ Class Components

- *Tactical components*
  - Stack
    - \* Used by non-recursive quick sort
  - Array
    - \* Stores/sorts pointers to lines & fields
  - Access\_Table
    - \* Used to store input
  - Input
    - \* Efficiently reads arbitrary sized input using only 1 dynamic allocation & 1 copy

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## C++ Class Components

- *Strategic components*
  - System\_Sort
    - \* Facade that integrates everything . . . .
  - Sort\_AT\_Adapter
    - \* Integrates Array & Access\_Table
  - Options
    - \* Manages globally visible options
  - Sort
    - \* e.g., both quicksort & insertion sort

---

## Detailed Format for Solution

- Note the separation of concerns

```
// Prototypes
template <typename ARRAY> void sort (ARRAY &a);
void operator<> (std::istream &, Sort_AT_Adapter &);
void operator<< (std::ostream &, const Sort_AT_Adapter &);

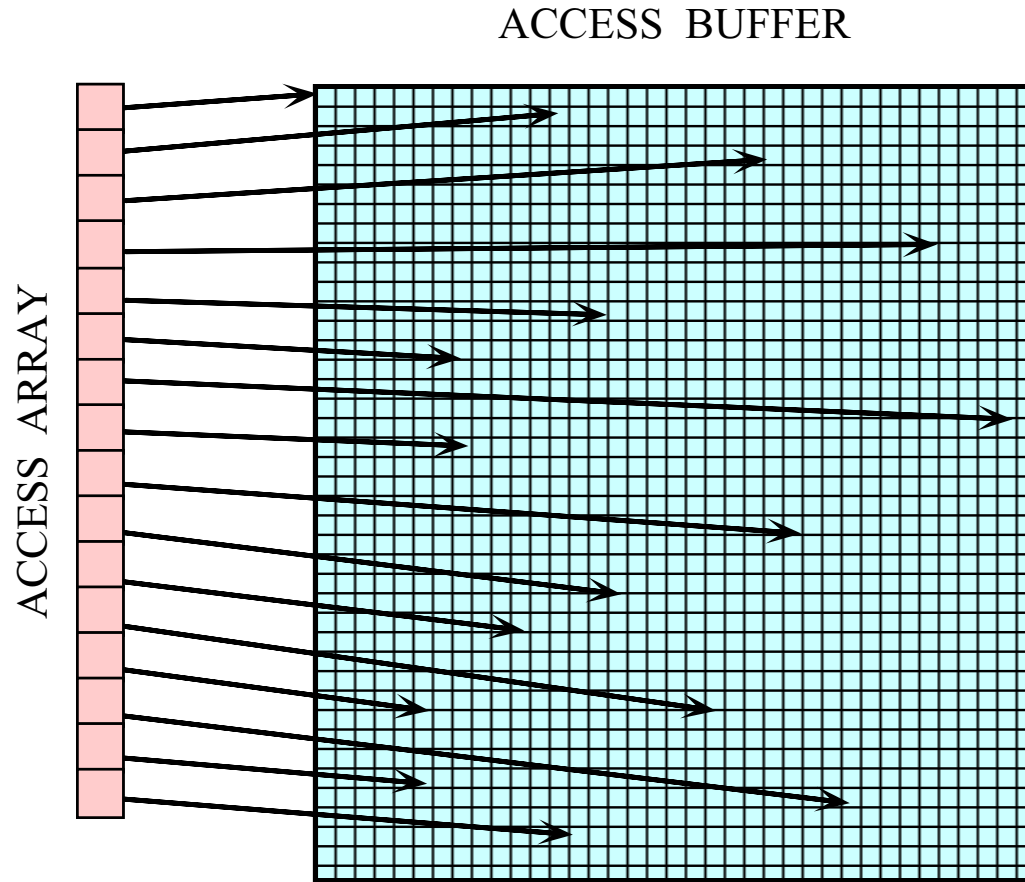
int main (int argc, char *argv[])
{
    Options::instance ()->parse_args (argc, argv);
    cin >> System_Sort::instance ()->access_table ();
    sort (System_Sort::instance ()->access_table ());
    cout << System_Sort::instance ()->access_table ();
}
```

---

## Reading Input Efficiently

- *Problem*
  - The input to the system sort can be arbitrarily large (e.g., up to 1/2 size of main memory)
- *Forces*
  - To improve performance solution must minimize:
    1. Data copying & data manipulation
    2. Dynamic memory allocation
- *Solution*
  - Create an Input class that reads arbitrary input efficiently

# Access Table Format



## The Input Class

- Efficiently reads arbitrary-sized input using only 1 dynamic allocation

```
class Input {
public:
    // Reads from <input> up to <terminator>, replacing <search>
    // with <replace>. Returns dynamically allocated buffer.
    char *read (std::istream &input, int terminator = EOF,
               int search = '\n', int replace = '\0');
    // Number of bytes replaced.
    size_t replaced () const;
    // Size of buffer.
    size_t size () const;
private:
    // Recursive helper method.
    char *recursive_read ();
    // . . .
};
```

## The Input Class (cont'd)

```
char *Input::read (std::istream &i, int t, int s, int r)
{
    // Initialize all the data members...
    return recursive_read ();
}

char *Input::recursive_read () {
    char buffer[BUFSIZ];
    // 1. Read input one character at a time, performing
    // search/replace until EOF is reached or buffer
    // is full.
    // 1.a If buffer is full, invoke recursive_read()
    // recursively.
    // 1.b If EOF is reached, dynamically allocate chunk
    // large enough to hold entire input
    // 2. On way out of recursion, copy buffer into chunk
}
```

---

## Design Patterns in the System Sort

- Facade
  - *Provide a unified interface to a set of interfaces in a subsystem*
    - \* Facade defines a higher-level interface that makes the subsystem easier to use
    - e.g., `sort()` function provides a facade for the complex internal details of efficient sorting
- Adapter
  - *Convert the interface of a class into another interface clients expect*
    - \* Adapter lets classes work together that couldn't otherwise because of incompatible interfaces
    - e.g., make `Access_Table` conform to interfaces expected by `sort & istd::ostreams`



---

## Design Patterns in System Sort (cont'd)

- Factory
  - Centralize assembly of resources needed to create objects
  - e.g., decouple initialization of Line\_Ptrs used by Access\_Table from their subsequent use
- Bridge
  - Decouple an abstraction from its implementation so that the two can vary independently
  - e.g., comparing two lines to determine ordering
- Strategy
  - Define a family of algorithms, encapsulate each one, & make them interchangeable
  - e.g., allow flexible pivot selection

---

## Design Patterns in System Sort (cont'd)

- Singleton
  - *Ensure a class has only one instance, & provide a global point of access to it*
  - e.g., provides a single point of access for the system sort facade & for program options
- Iterator
  - *Provide a way to access the elements of an aggregate object sequentially without exposing its underlying representation*
  - e.g., provides a way to print out the sorted lines without exposing representation or initialization

---

## Sort Algorithm

- For efficiency, two types of sorting algorithms are used:
  1. *Quicksort*
    - Highly time & space efficient sorting arbitrary data
    - $O(n \log n)$  average-case time complexity
    - $O(n^2)$  worst-case time complexity
    - $O(\log n)$  space complexity
    - Optimizations are used to avoid worst-case behavior
  2. *Insertion sort*
    - Highly time & space efficient for sorting “almost ordered” data
    - $O(n^2)$  average- & worst-case time complexity
    - $O(1)$  space complexity

---

## Quicksort Optimizations

1. *Non-recursive*
  - Uses an explicit stack to reduce function call overhead
2. *Median of 3 pivot selection*
  - Reduces probability of worse-case time complexity
3. *Guaranteed ( $\log n$ ) space complexity*
  - Always “pushes” larger partition
4. *Insertion sort for small partitions*
  - Insertion sort runs fast on almost sorted data

---

## Selecting a Pivot Value

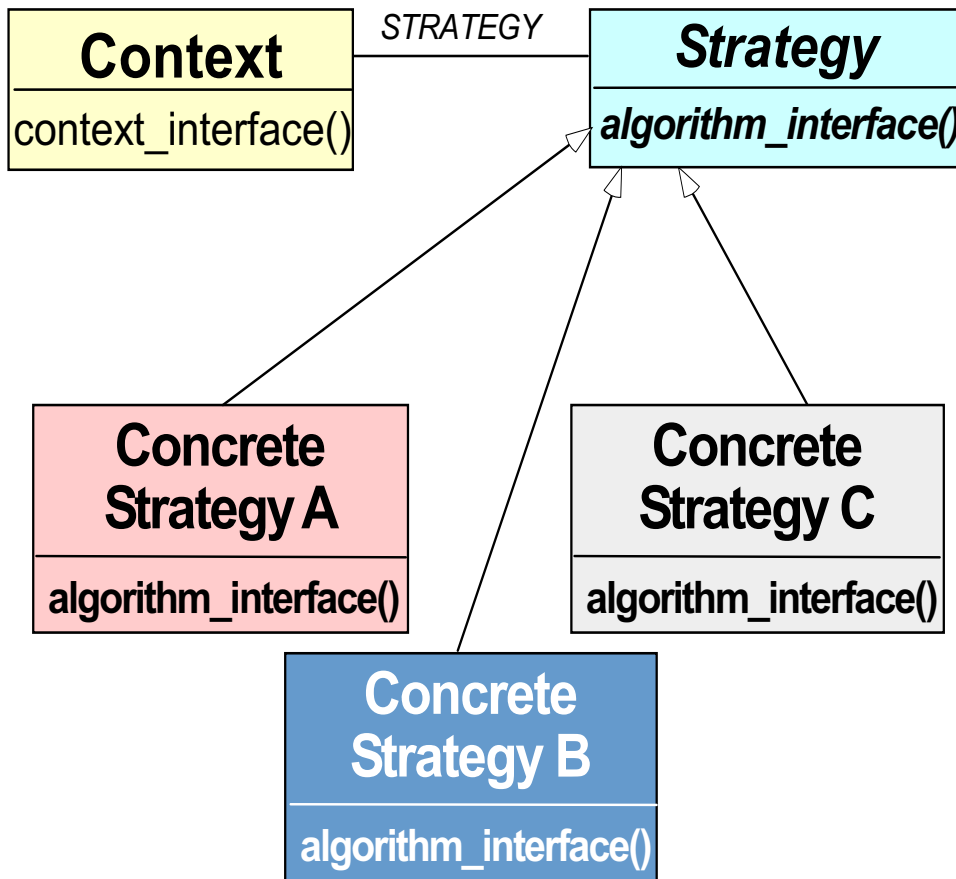
- *Problem*
  - There are various algorithms for selecting a pivot value
    - \* e.g., randomization, median of three, etc.
- *Forces*
  - Different input may sort more efficiently using different pivot selection algorithms
- *Solution*
  - Use the *Strategy* pattern to select the pivot selection algorithm

---

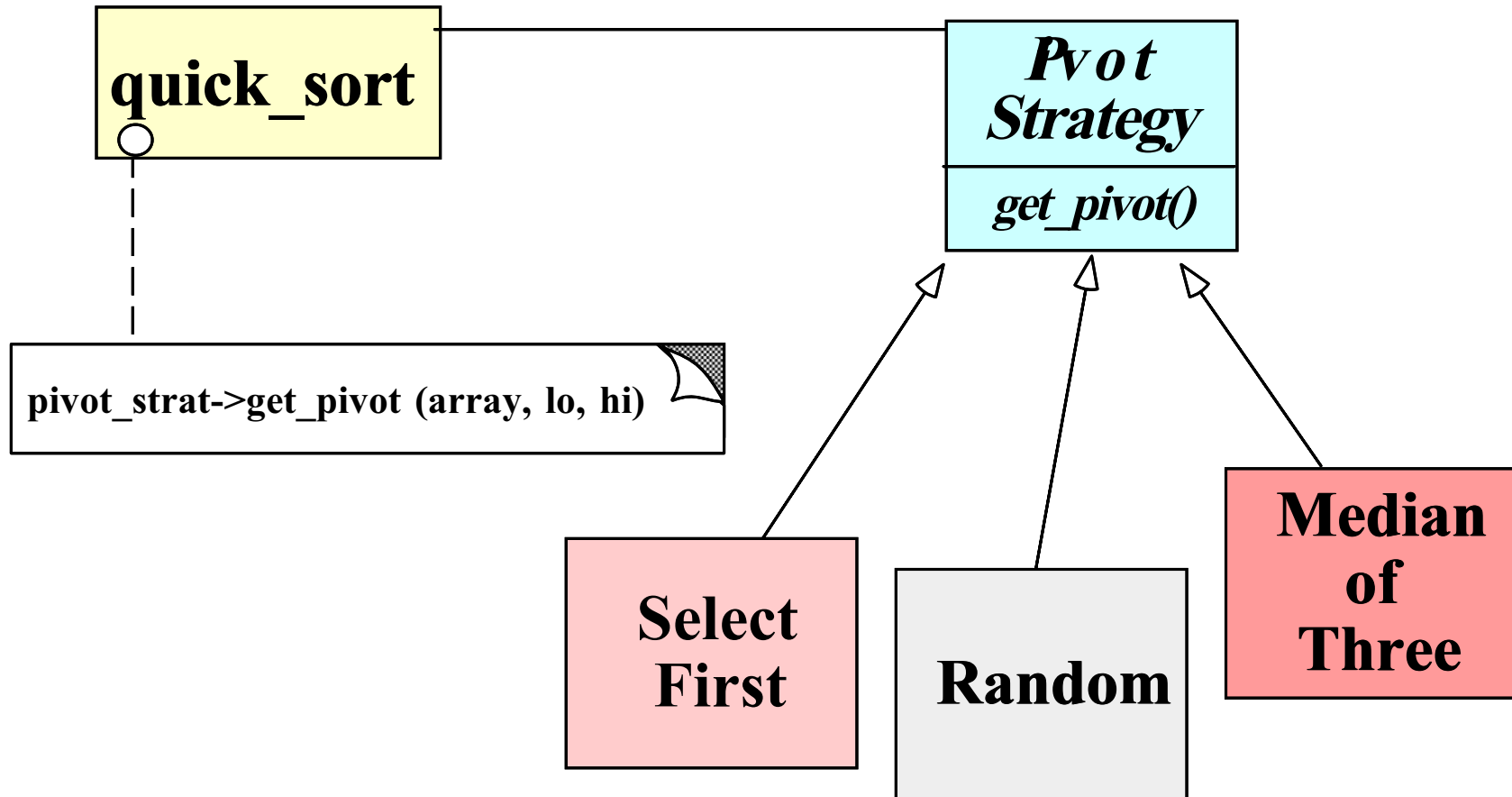
## The Strategy Pattern

- *Intent*
  - Define a family of algorithms, encapsulate each one, & make them interchangeable
    - \* Strategy lets the algorithm vary independently from clients that use it
- This pattern resolves the following forces
  1. *How to extend the policies for selecting a pivot value without modifying the main quicksort algorithm*
  2. *Provide a one size fits all interface without forcing a one size fits all implementation*

# Structure of the Strategy Pattern



# Using the Strategy Pattern





## Implementing the Strategy Pattern

- ARRAY is the particular “context”

```
template <typename ARRAY>
void sort (ARRAY &array) {
    Pivot_Strategy<ARRAY> *pivot_strat =
        Pivot_Factory<ARRAY>::make_pivot
            (Options::instance ()->pivot_strat ());
    std::auto_ptr <Pivot_Strategy<ARRAY> >
        holder (pivot_strat);

    // Ensure exception safety.
    ARRAY temp = array;
    quick_sort (temp, pivot_strat);
    // Destructor of <holder> deletes <pivot_strat>.
    array = temp;
}
```

## Implementing the Strategy Pattern

```
template <typename ARRAY, class PIVOT_STRAT>
quick_sort (ARRAY &array,
            PIVOT_STRAT *pivot_strat) {
    for (;;) {
        typename ARRAY::TYPE pivot =
            // Note 'lo' & 'hi' should be passed by reference
            // so get_pivot() can reorder the values & update
            // 'lo' & 'hi' accordingly...
            pivot_strat->get_pivot (array, lo, hi);
        // Partition array[lo, hi] relative to pivot . . .
    }
}
```

## Fixed-size Stack

- Defines a fixed size stack for use with non-recursive quicksort

```
template <typename T, size_t SIZE>
class Fixed_Stack
{
public:
    bool push (const T &new_item);
    bool pop (T &item);
    bool is_empty ();
    // . . .

private:
    T stack_[SIZE];
    size_t top_;
};
```

---

## Devising a Simple Sort Interface

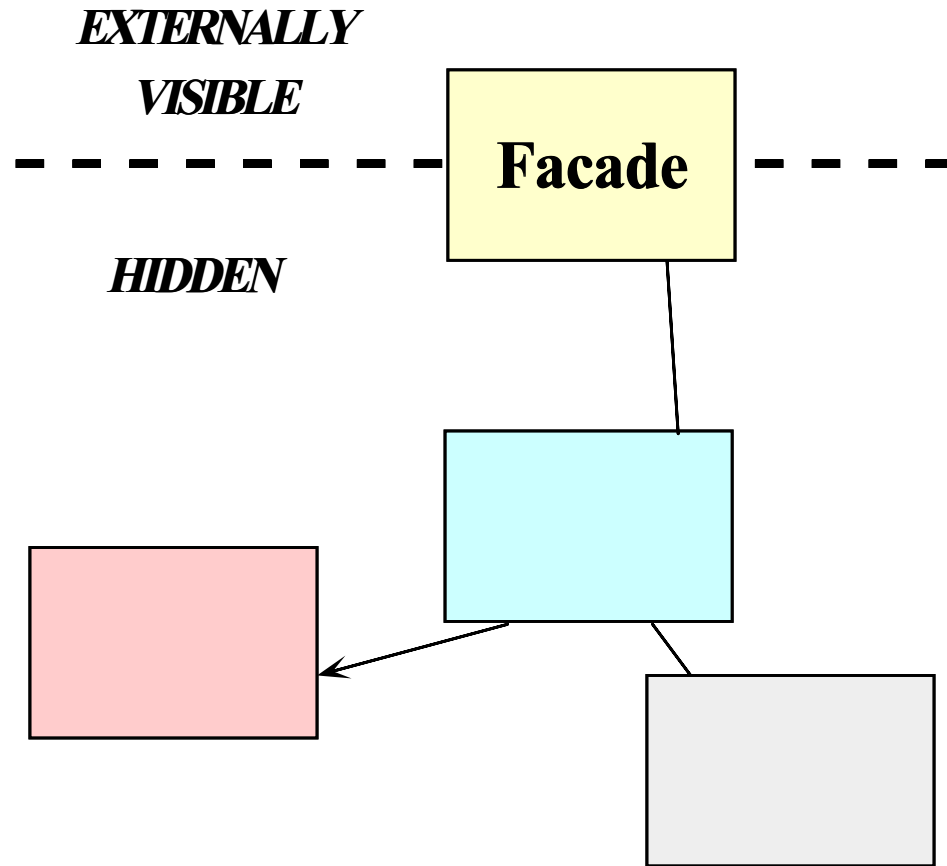
- *Problem*
  - Although the implementation of the sort function is complex, the interface should be simple to use
- *Key forces*
  - Complex interface are hard to use, error prone, and discourage extensibility & reuse
  - Conceptually, sorting only makes a few assumptions about the “array” it sorts
    - \* e.g., supports operator[] methods, size, & trait TYPE
  - We don’t want to arbitrarily limit types of arrays we can sort
- *Solution*
  - Use the *Facade* & *Adapter* patterns to simplify the sort program

---

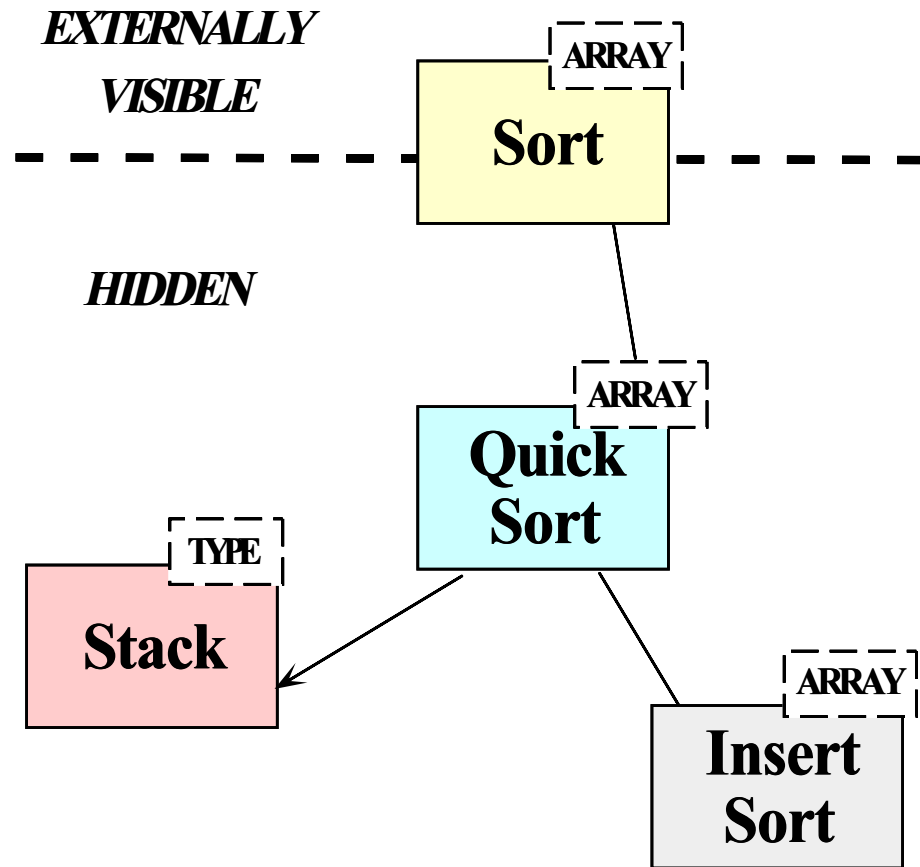
## Facade Pattern

- *Intent*
  - Provide a unified interface to a set of interfaces in a subsystem
    - \* Facade defines a higher-level interface that makes the subsystem easier to use
- This pattern resolves the following forces:
  1. Simplifies the sort interface
    - e.g., only need to support operator [] & size methods, & element TYPE
  2. Allows the implementation to be efficient and arbitrarily complex without affecting clients

# Structure of the Facade Pattern



# Using the Facade Pattern



---

## Centralizing Option Processing

- *Problem*
  - Command-line options must be global to many parts of the sort program
- *Key forces*
  - Unrestricted use of global variables increases system coupling & can violate encapsulation
  - Initialization of static objects in C++ can be problematic
- *Solution*
  - Use the *Singleton* pattern to centralize option processing



---

## Singleton Pattern

- *Intent*
  - *Ensure a class has only one instance, & provide a global point of access to it*
- This pattern resolves the following forces:
  1. Localizes the creation & use of “global” variables to well-defined objects
  2. Preserves encapsulation
  3. Ensures initialization is done after program has started & only on first use
  4. Allow transparent subclassing of Singleton implementation

## Structure of the Singleton Pattern

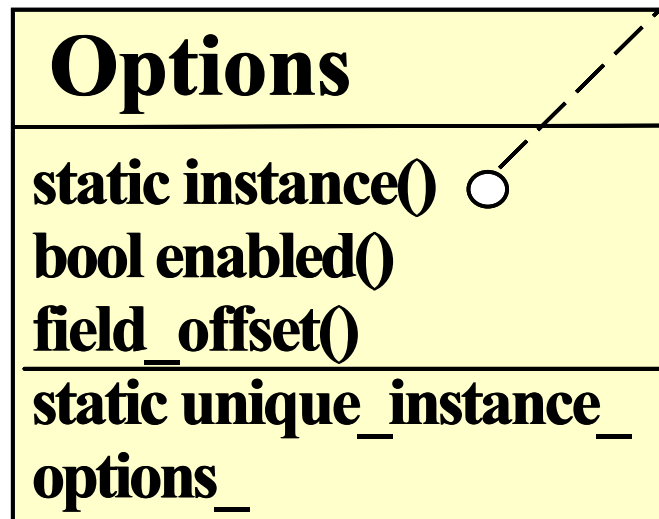
```
if (unique_instance_ == 0)
    unique_instance_ = new Singleton;
return unique_instance_;
```

### Singleton

```
static instance() ○
singleton_operation()
get_singleton_data()
static unique_instance_
singleton_data_
```

## Using the Singleton Pattern

```
if (unique_instance_ == 0)
    unique_instance_ = new Options;
return unique_instance_;
```



## Options Class

- This manages globally visible options

```
class Options
{
public:
    static Options *instance ();
    bool parse_args (int argc, char *argv []);

    // These options are stored in octal order
    // so that we can use them as bitmasks!
    enum Option { FOLD = 01, NUMERIC = 02,
                 REVERSE = 04, NORMAL = 010 };
    enum Pivot_Strategy { MEDIAN, RANDOM, FIRST };
};
```

## Options Class (cont'd)

```
bool enabled (Option o);

int field_offset (); // Offset from BOL.
Pivot_Strategy pivot_strat ();
int (*compare) (const char *l, const char *r);

protected:
Options (); // Ensure Singleton.

u_long options_; // Maintains options bitmask . . .
int field_offset_;
static Options *instance_; // Singleton.
};
```

## Options Class (cont'd)

```
#define SET_BIT(WORD, OPTION) (WORD |= OPTION)
#define CLR_BIT(WORD, OPTION) (WORD &= ~OPTION)

bool Options::parse_args (int argc, char *argv[]) {
    for (int c;
         (c = getopt (argc, argv, 'nrfs:k:c:t:')) != EOF; ) {
        switch (c) {
            case 'n': {
                CLR_BIT (options_, Options::FOLD);
                CLR_BIT (options_, Options::NORMAL);
                SET_BIT (options_, Options::NUMERIC);
                break;
            }
            // . . .
        }
    }
}
```

## Using the Options Class

- One way to implement `sort()` comparison operator:

```
int Line_Ptrs::operator< (const Line_Ptrs &rhs) const {
    Options *options = Options::instance ();
    if (options->enabled (Options::NORMAL))
        return strcmp (this->bof_, rhs.bof_) < 0;
    else if (options->enabled (Options::NUMERIC));
        return numcmp (this->bof_, rhs.bof_) < 0;
    else // if (options->enabled (Options::FOLD))
        return strcasecmp (this->bof_, rhs.bof_) < 0;
}
```

- We'll see another approach later on using Bridge

---

## Simplifying Comparisons

- *Problem*
  - The comparison operator shown above is somewhat complex
- *Forces*
  - It's better to determine the type of comparison operation during the initialization phase
  - But the interface shouldn't change
- *Solution*
  - Use the *Bridge pattern* to separate interface from implementation

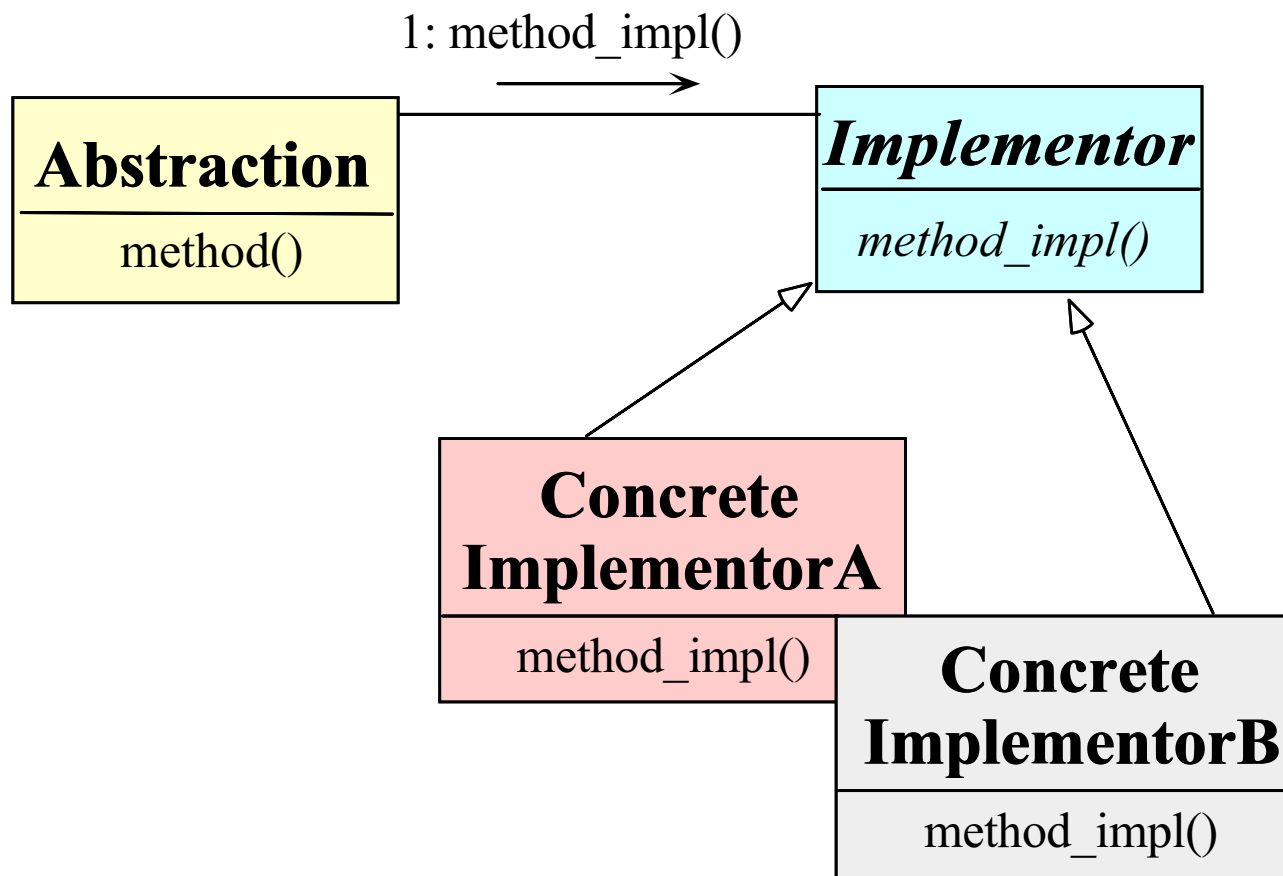


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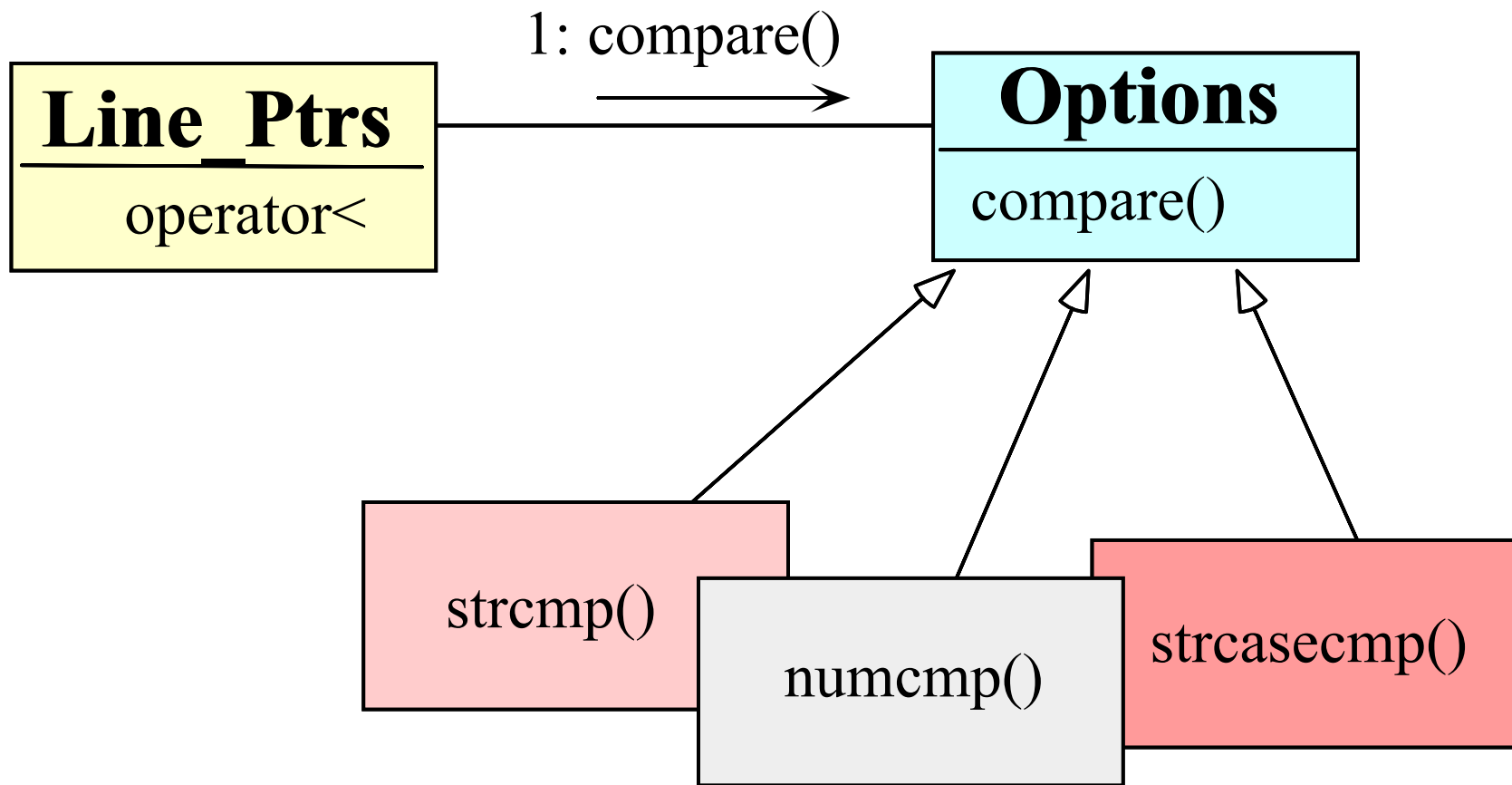
## The Bridge Pattern

- *Intent*
  - *Decouple an abstraction from its implementation so that the two can vary independently*
- This pattern resolves the following forces that arise when building extensible software
  1. *How to provide a stable, uniform interface that is both closed & open, i.e.,*
    - *Closed to prevent direct code changes*
    - *Open to allow extensibility*
  2. *How to simplify the `Line_Ptrs::operator<` implementation & reference counting for `Access_Table` buffer*

## Structure of the Bridge Pattern



# Using the Bridge Pattern



## Using the Bridge Pattern

- The following is the comparison operator used by sort

```
int Line_Ptrs::operator<(const Line_Ptrs &rhs) const {  
    return (*Options::instance()->compare)  
        (bof_, rhs.bof_) < 0;  
}
```

- This solution is much more concise
- However, there's an extra level of function call indirection . . .
  - Which is equivalent to a virtual function call

---

## Initializing the Comparison Operator

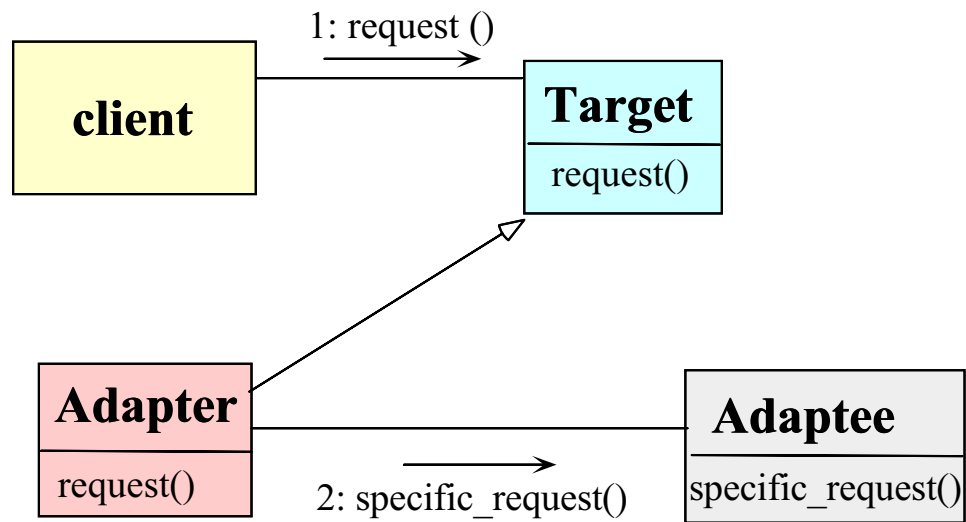
- *Problem*
  - How does the compare pointer-to-method get assigned?  
`int (*compare) (const char *left, const char *right);`
- *Forces*
  - There are many different choices for compare, depending on which options are enabled
  - We only want to worry about initialization details in one place
  - Initialization details may change over time
  - We'd like to do as much work up front to reduce overhead later on
- *Solution*
  - Use a *Factory* pattern to initialize the comparison operator

---

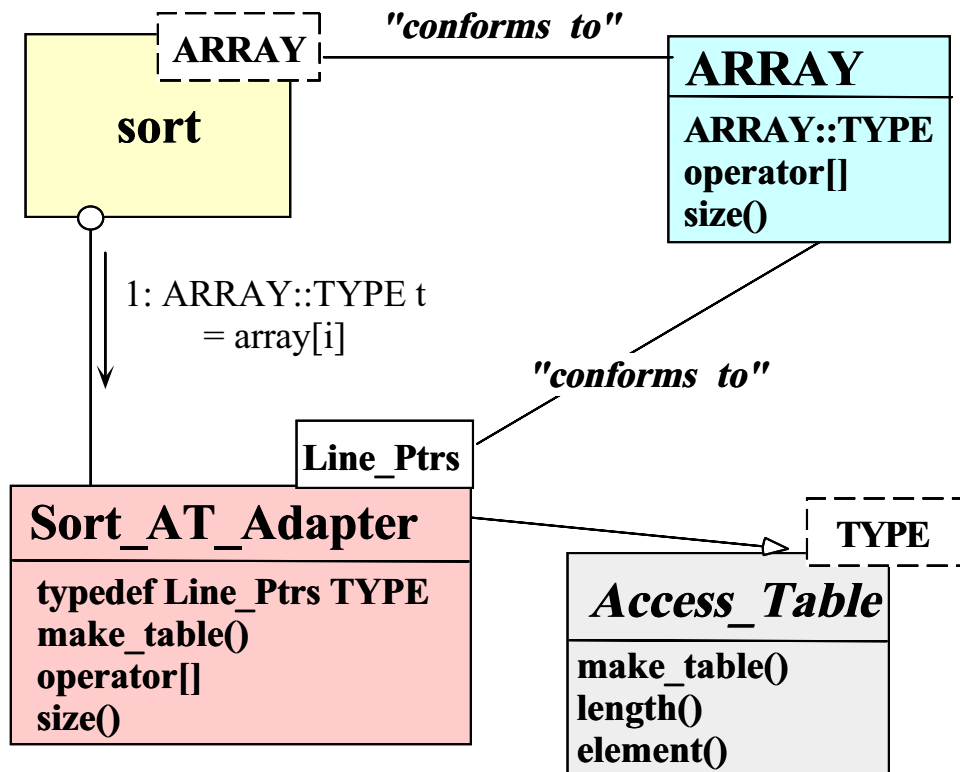
## The Adapter Pattern

- *Intent*
  - *Convert the interface of a class into another interface clients expect*
    - \* Adapter lets classes work together that couldn't otherwise because of incompatible interfaces
- This pattern resolves the following forces:
  1. How to transparently integrate the `Access_Table` with the `sort` routine
  2. How to transparently integrate the `Access_Table` with the C++ `istd::ostream` operators

## Structure of the Adapter Pattern



# Using the Adapter Pattern





## Dynamic Array

- Defines a variable-sized array for use by the Access\_Table

```
template <typename T>
class Array {
public:
    Array (size_t size = 0);
    int init (size_t size);
    T &operator [] (size_t index);
    size_t size () const;
    T *begin () const; // STL iterator methods.
    T *end () const;
    // . . .
private:
    T *array_;
    size_t size_;
};
```



## The Access\_Table Class

- Efficiently maps indices onto elements in the data buffer

```
template <typename T>
class Access_Table {
public:
    // Factory Method for initializing Access_Table.
    virtual int make_table (size_t lines, char *buffer) = 0;
    // Release buffer memory.
    virtual ~Access_Table ();
    T &element (size_t index); // Reference to <indexth> element.
    size_t length () const; // Length of the access_array.
    Array<T> &array (void) const; // Return reference to array.
protected:
    Array<T> access_array_; // Access table is array of T.
    Access_Table_Impl *access_table_impl_; // Ref counted buffer.
};
```

## The Access\_Table\_Impl Class

```
class Access_Table_Impl { // Part of the Bridge pattern
public:
    Access_Table_Impl (void); //Default constructor
    Access_Table_Impl (char *buffer); // Constructor
    // Virtual destructor ensures subclasses are virtual
    virtual ~Access_Table_Impl (void);

    void add_ref (void); // Increment reference count
    void remove_ref (void); // Decrement reference count
    char *get_buffer(void); // Get buffer from the class
    void set_buffer(char *); // Set buffer

private:
    char *buffer_; // Underlying buffer
    size_t ref_count_; // Refcount tracks deletion.
};
```

---

## The Sort\_AT\_Adapter Class

- Adapts the Access\_Table to conform to the ARRAY interface expected by

```
sort
```

```
struct Line_Ptrs {  
    // Comparison operator used by sort().  
    int operator< (const Line_Ptrs &) const;  
  
    // Beginning of line & field/column.  
    char *bol_, *bof_;  
};
```

## The Sort\_AT\_Adapter Class

```
class Sort_AT_Adapter : // Note class form of the Adapter
private Access_Table<Line_Ptrs> {
public:
    virtual int make_table (size_t num_lines, char *buffer);

    typedef Line_Ptrs TYPE; // Type trait.

    // These methods adapt Access_Table methods . . .
    Line_Ptrs &operator[] (size_t index);
    size_t size () const;
};

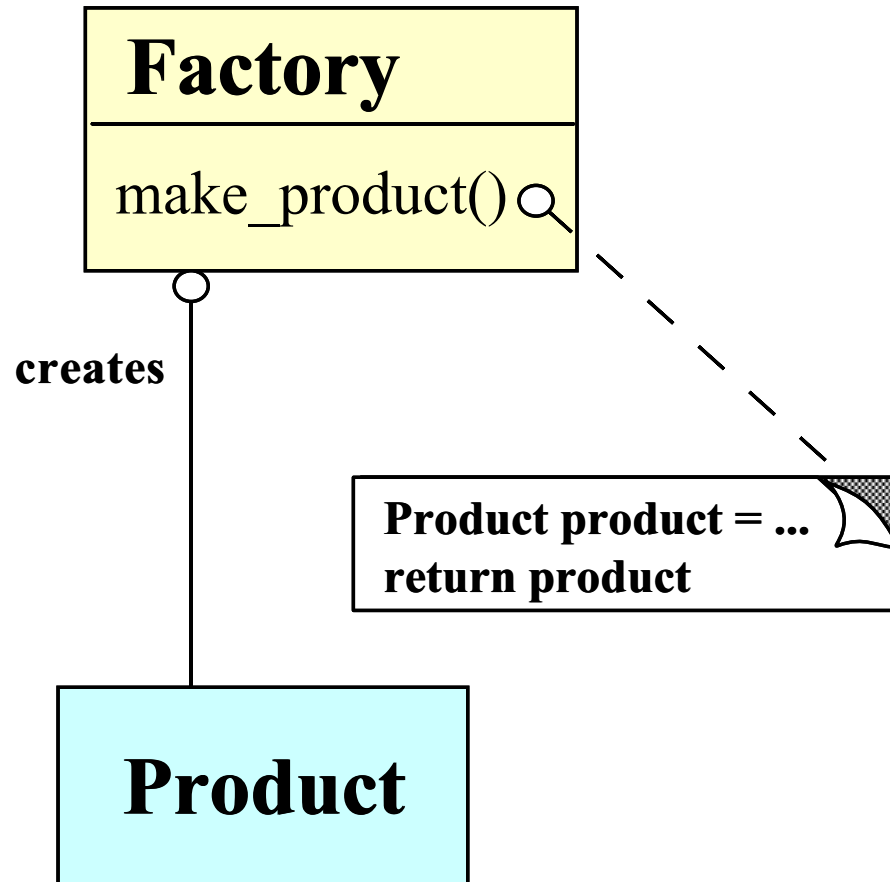
// Put these into separate file.
Line_Ptrs &Sort_AT_Adapter::operator[] (size_t i)
{ return element (i); }
size_t Sort_AT_Adapter::size () const { return length (); }
```

---

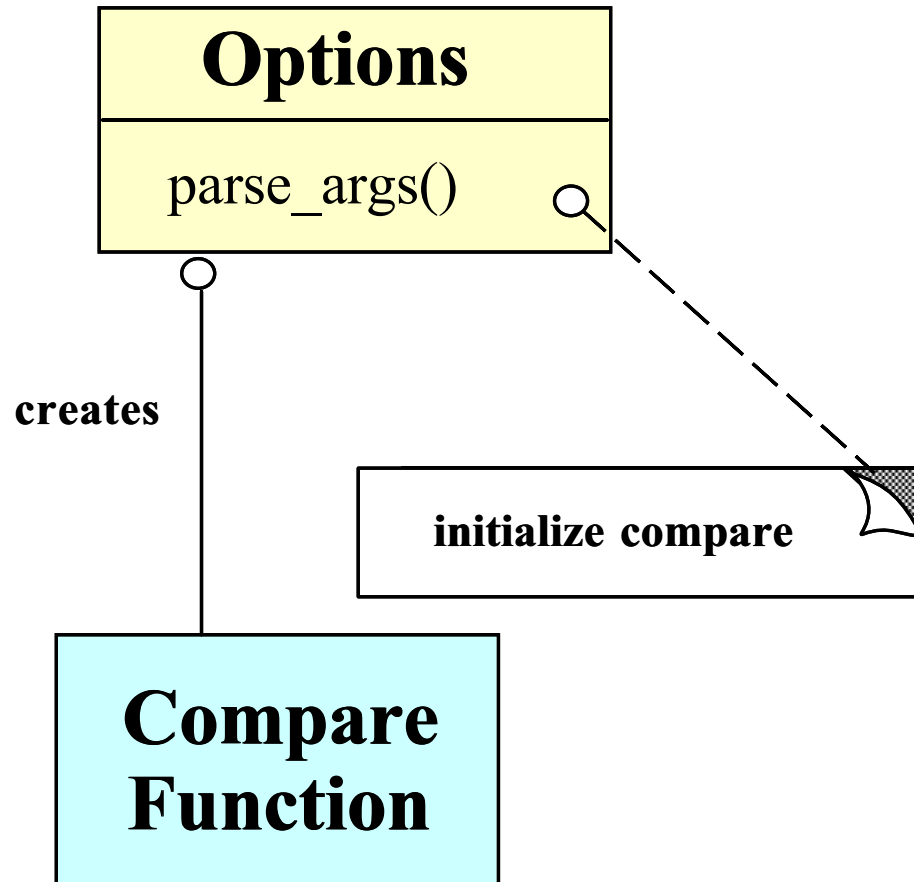
## The Factory Pattern

- *Intent*
  - *Centralize the assembly of resources necessary to create an object*
    - \* Decouple object creation from object use by localizing creation knowledge
- This pattern resolves the following forces:
  - Decouple initialization of the compare operator from its subsequent use
  - Makes it easier to change comparison policies later on
    - \* e.g., adding new command-line options

## Structure of the Factory Pattern



# Using the Factory Pattern for Comparisons





## Code for Using the Factory Pattern

- The following initialization is done after command-line options are parsed

```
bool Options::parse_args (int argc, char *argv[])
{
    // . . .
    if (this->enabled (Options::NORMAL))
        this->compare = &strcmp;
    else if (this->enabled (Options::NUMERIC))
        this->compare = &numcmp;
    else if (this->enabled (Options::FOLD))
        this->compare = &strcasecmp;
    // . . .
}
```

---

## Code for Using the Factory Pattern (cont'd)

- We need to write a `numcmp()` adapter function to conform to the API used by the compare pointer-to-function

```
int numcmp (const char *s1, const char * s2) {  
    double d1 = strtod (s1, 0), d2 = strtod (s2, 0);  
  
    if (d1 < d2) return -1;  
    else if (d1 > d2) return 1;  
    else // if (d1 == d2)  
        return 0;  
}
```

---

## Initializing the Access\_Table

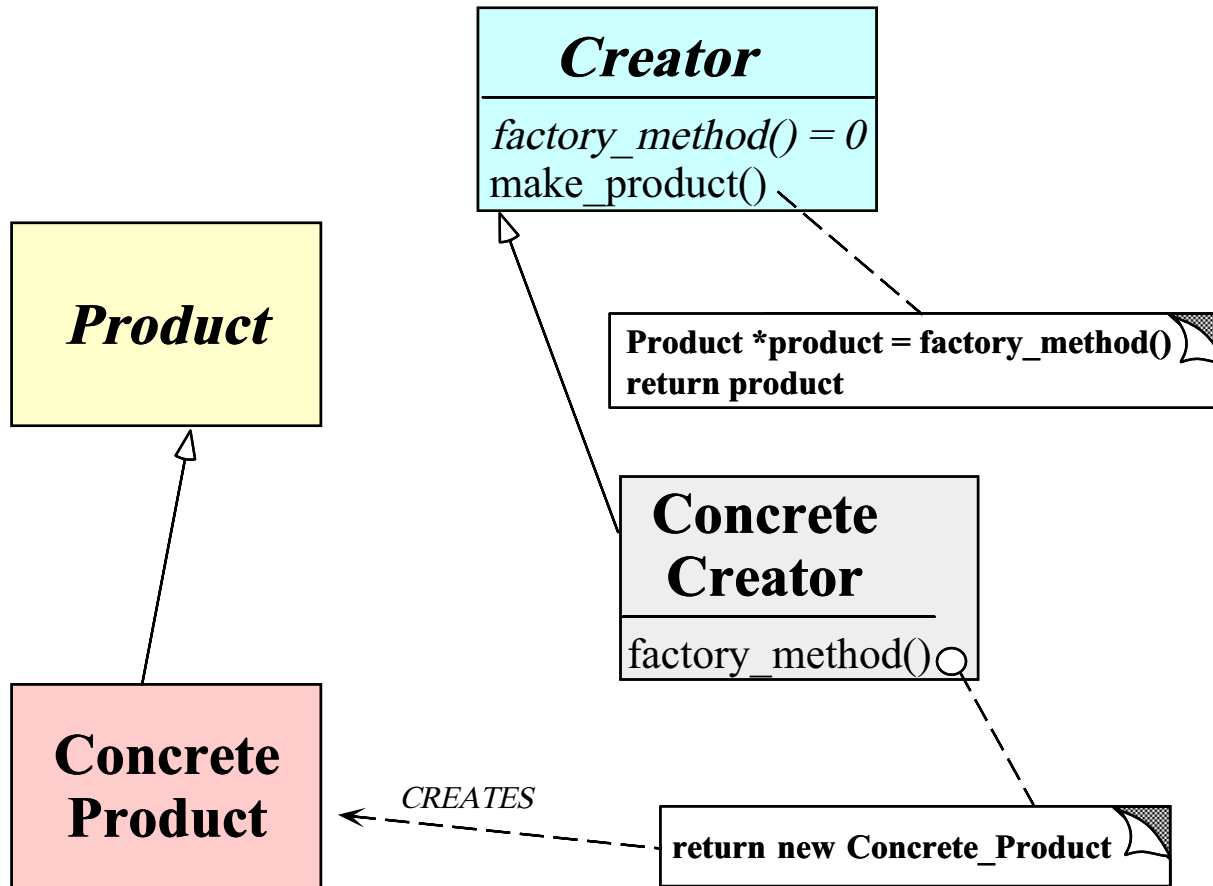
- *Problem*
  - One of the nastiest parts of the whole system sort program is initializing the Access\_Table
- *Key forces*
  - We don't want initialization details to affect subsequent processing
  - Makes it easier to change initialization policies later on
    - \* e.g., using the Access\_Table in non-sort applications
- *Solution*
  - Use the *Factory Method* pattern to initialize the Access\_Table

---

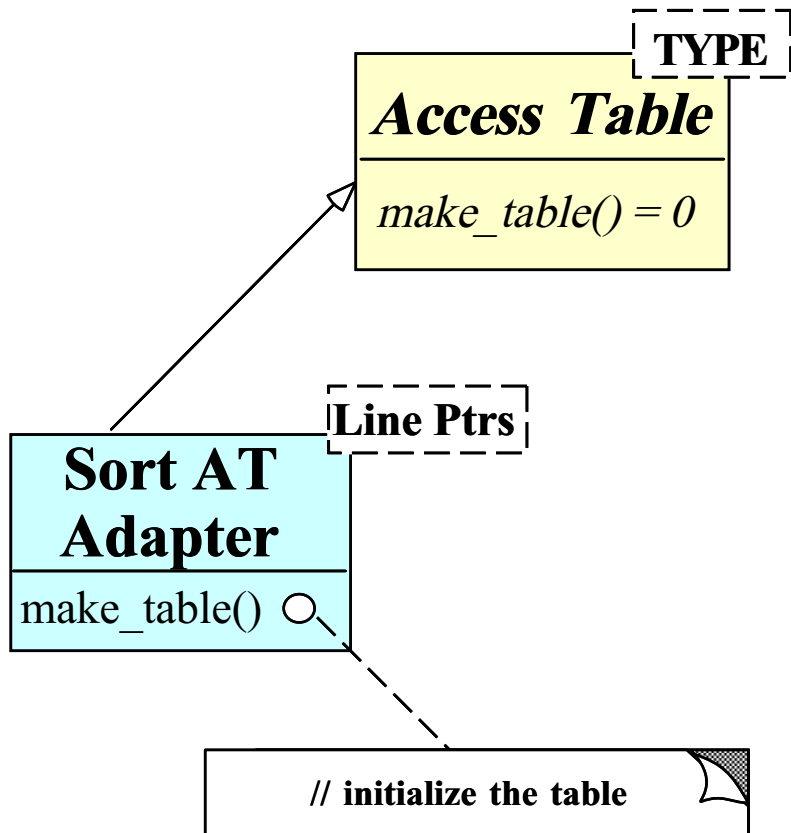
## Factory Method Pattern

- *Intent*
  - Define an interface for creating an object, but let subclasses decide which class to instantiate
    - \* Factory Method lets a class defer instantiation to subclasses
- This pattern resolves the following forces:
  - Decouple initialization of the `Access_Table` from its subsequent use
  - Improves subsequent performance by pre-caching beginning of each field & line
  - Makes it easier to change initialization policies later on
    - \* e.g., adding new command-line options

# Structure of the Factory Method Pattern



# Using the Factory Method Pattern for Access\_Table Initialization



## Using the Factory Method Pattern for the Sort\_AT\_Adapter

- The following istd::ostream Adapter initializes the Sort\_AT\_Adapter access table

```
void operator>> (std::istream &is, Sort_AT_Adapter &at)
{
    Input input;
    // Read entire stdin into buffer.
    char *buffer = input.read (is);
    size_t num_lines = input.replaced ();

    // Factory Method initializes Access_Table<>.
    at.make_table (num_lines, buffer);
}
```

---

## Implementing the Factory Method Pattern

- The Access\_Table\_Factory class has a Factory Method that initializes Sort\_AT\_Adapter

```
// Factory Method initializes Access_Table.  
int Sort_AT_Adapter::make_table (size_t num_lines,  
                                char *buffer)  
{  
    // Array assignment op.  
    this->access_array_.resize (num_lines);  
    this->buffer_ = buffer; // Obtain ownership.  
  
    size_t count = 0;
```



## Implementing the Factory Method Pattern (cont'd)

```
// Iterate through the buffer & determine
// where the beginning of lines & fields
// must go.
for (Line_Ptrs_Iter iter (buffer, num_lines);
     iter.is_done () == 0;
     iter.next ())
{
    Line_Ptrs line_ptr = iter.current_element ();
    this->access_array_[count++] = line_ptr;
}
}
```

---

## Initializing the Access\_Table with Input Buffer

- *Problem*
  - We'd like to initialize the `Access_Table` *without* having to know the input buffer is represented
- *Key force*
  - Representation details can often be decoupled from accessing each item in a container or collection
- *Solution*
  - Use the *Iterator* pattern to scan through the buffer

---

## Iterator Pattern

- *Intent*
  - Provide a way to access the elements of an aggregate object sequentially without exposing its underlying representation
- The C++ Standard Library (STL) is heavily based on the iterator pattern, e.g.,

```
int main (int argc, char *argv[]) {
    std::vector <std::string> args;
    for (int i = 1; i < argc; ++i) {
        args.push_back (std::string (argv [i]));
    }
    for (std::vector<std::string>::iterator j = args.begin ();
        j != args.end (); ++j)
        cout << (*j)_ << endl;
}
```

## Iterator Pattern (cont'd)

- The Iterator pattern provides a way to initialize the Access\_Table without knowing how the buffer is represented

```
Line_Ptrs_Iter::Line_Ptrs_Iter (char *buffer,  
                               size_t num_lines);
```

```
Line_Ptrs Line_Ptrs_Iter::current_element () {  
    Line_Ptrs lp;
```

```
    // Determine beginning of next line \& next field . . .  
    lp.bof_ = // . . .  
    lp.bof_ = // . . .
```

```
    return lp;  
}
```

## Iterator Pattern (cont'd)

- Iterator provides a way to print out sorted lines

```
void operator<< (std::ostream &os, const Line_Ptrs lp) {
    os << lp.bol_;
}

void operator<< (std::ostream &os, const Sort_AT_Adapter &at) {
    if (Options::instance ()->enabled (Options::REVERSE))
        std::reverse_copy (
            at.array ().begin (),
            at.array ().end (),
            std::ostream_iterator<System_Sort::Line_Ptrs> (os, "\n"));
    else
        std::copy (
            at.array ().begin (),
            at.array ().end (),
            std::ostream_iterator<System_Sort::Line_Ptrs> (os, "\n"));
}
```

---

## Summary of System Sort Case Study

- This case study illustrates using OO techniques to structure a modular, reusable, & highly efficient system
- Design patterns help to resolve many key forces
- Performance of our system sort is comparable to existing UNIX system sort
  - Use of C++ features like *parameterized types* and *inlining* minimizes penalty from increased modularity, abstraction, & extensibility

---

## Case Study: Sort Verifier

- *Verify whether a sort routine works correctly*
  - *i.e.*, output of the sort routine must be an ordered permutation of the original input
- This is useful for checking our system sort routine!
- The solution is harder than it looks at first glance . . .
- As before, we'll examine the key forces & discuss design patterns that resolve the forces

## General Form of Solution

- The following is a general use-case for this routine:

```
template <typename ARRAY> void sort (ARRAY &a);
```

```
template <typename ARRAY> int
```

```
check_sort (const ARRAY &o, const ARRAY &p);
```

```
int main (int argc, char *argv[])  
{
```

```
    Options::instance ()->parse_args (argc, argv);
```

```
    Input original;
```

```
    Input potentially_sorted;
```



## General Form of Solution (cont'd)

```
cin >> input;

std::copy (original.begin (),
          original.end (),
          potentially_sorted.begin ());
sort (potentially_sorted);

if (check_sort (original, potentially_sorted) == -1)
    cerr << "sort failed" << endl;
else
    cout << "sort worked" << endl;
}
```

## Common Problems

unsorted	7	13	4	15	18	13	8	4
sorted, but not permuted	0	0	0	0	0	0	0	0
permuted, but not sorted	8	13	18	15	4	13	4	7
sorted and permuted	4	4	7	8	13	13	15	18

- Several common problems:
  - Sort routine may zero out data
    - \* though it will appear sorted . . . ;-)
  - Sort routine may fail to sort data
  - Sort routine may erroneously add new values

---

## Forces

- Solution should be both time & space efficient
  - e.g., it should not take more time to check than to sort in the first place!
  - Also, this routine may be run many times consecutively, which may facilitate certain space optimizations
- We cannot assume the existence of a “correct” sorting algorithm . . .
  - Therefore, to improve the chance that our solution is correct, it must be simpler than writing a correct sorting routine
    - \* *Quis custodiet ipsos custodes?*
      - (Who shall guard the guardians?)

## Forces (cont'd)

- Multiple implementations will be necessary, depending on properties of the data being examined, e.g.,
  1. if data values are small (in relation to number of items) & integrals use . . . .
  2. if data has no duplicate values use . . . .
  3. if data has duplicate values use . . . .
- This problem illustrates a simple example of “program families”
  - *i.e.*, we want to reuse as much code and/or design across multiple solutions as possible

---

## Strategies

- Implementations of search structure vary according to data, e.g.,
  1. *Range Vector*
    - $O(N)$  time complexity & space efficient for sorting “small” ranges of integral values
  2. *Binary Search (version 1)*
    - $O(n \log n)$  time complexity & space efficient but does not handle duplicates
  3. *Binary Search (version 2)*
    - $O(n \log n)$  time complexity, but handles duplicates
  4. *Hashing*
    - $O(n)$  best/average case, but  $O(n^2)$  worst case, handles duplicates, but potentially not as space efficient

---

## General OOD Solution Approach

- Identify the “objects” in the application & solution space
  - e.g., use a *search structure ADT* organization with member function such as *insert & remove*
- Recognize common design patterns
  - e.g., Strategy & Factory Method
- Implement a framework to coordinate multiple implementations
  - e.g., use classes, parameterized types, inheritance & dynamic binding

---

## General OOD solution approach (cont'd)

- C++ framework should be amenable to:
  - *Extension & Contraction*
    - \* May discover better implementations
    - \* May need to conform to resource constraints
    - \* May need to work on multiple types of data
  - *Performance Enhancement*
    - \* May discover better ways to allocate & cache memory
    - \* Note, improvements should be transparent to existing code . . .
  - *Portability*
    - \* May need to run on multiple platforms

## High-level Algorithm

- e.g., pseudo code

```
template <typename ARRAY>
int check_sort (const ARRAY &original,
               const ARRAY &potential_sort)
{
```

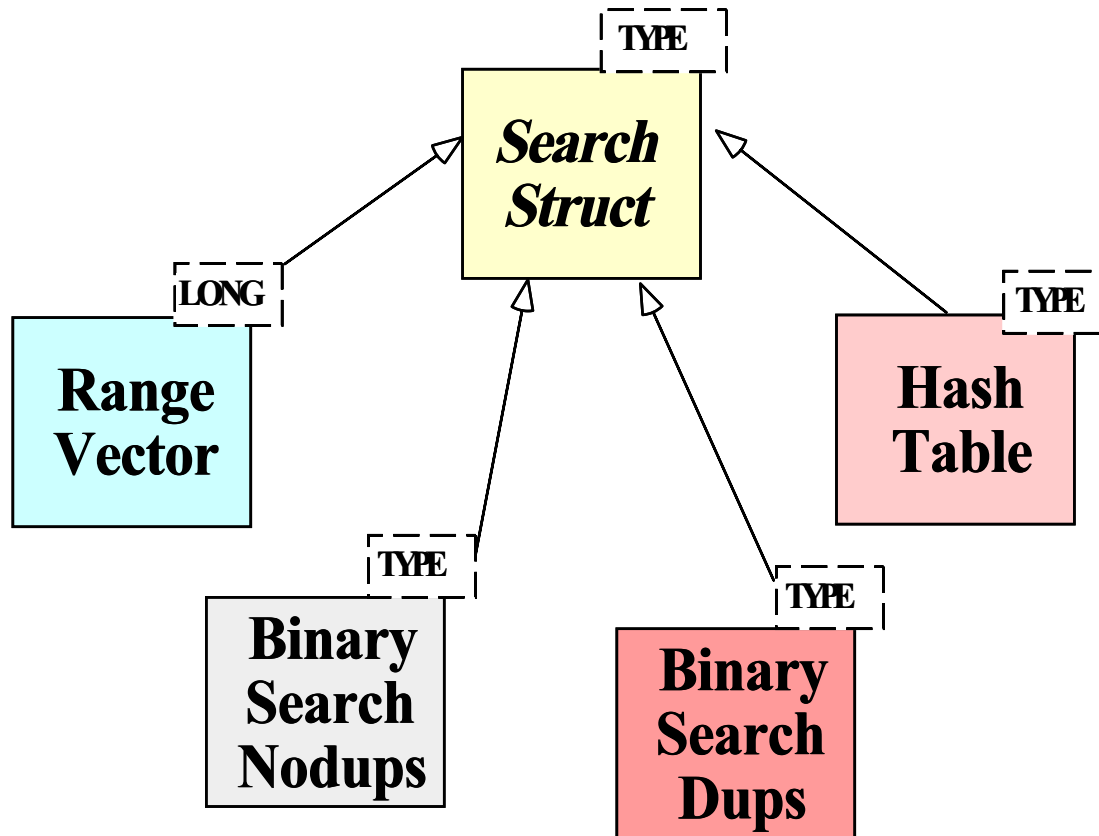
    Perform basic sanity check to see if the  
    potential\_sort is actually in order  
    (can also detect duplicates here)



## High-level Algorithm (cont'd)

```
if (basic sanity check succeeds) then
  Initialize search structure, srchstrct
  for i < 0 to size - 1 loop
    insert (potential_sort[i])
      into srchstrct
  for i < 0 to size - 1 loop
    if remove (original[i]) from
      srchstrct fails then
      return ERROR
    return SUCCESS
  else
    return ERROR
  end if
}
```

## UML Class Diagram for C++ Solution



## C++ Class Interfaces

- Search structure base class.

```
template <typename T>
class Search_Strategy
{
public:
    virtual bool insert (const T &new_item) = 0;
    virtual bool remove (const T &existing_item) = 0;
    virtual ~Search_Strategy () = 0;
};
```

## C++ Class interfaces (cont'd)

- Strategy Factory class

```
template <typename ARRAY>
Search_Struct
{
public:
    // Singleton method.
    static Search_Struct<ARRAY> *instance ();

    // Factory Method
    virtual Search_Strategy<typename ARRAY::TYPE> *
        make_strategy (const ARRAY &);
};
```

## C++ Class interfaces (cont'd)

- Strategy subclasses

```
// Note the template specialization
class Range_Vector :
    public Search_Strategy<long>
{ typedef long TYPE; /* . . . */ };

template <typename ARRAY>
class Binary_Search_Nodups :
    public Search_Strategy<typename ARRAY::TYPE>
{
    typedef typename ARRAY::TYPE TYPE; /* . . . */
};
```

## C++ Class interfaces (cont'd)

```
template <typename ARRAY> class Binary_Search_Dups :
public Search_Strategy<typename ARRAY::TYPE>
{
    typedef typename ARRAY::TYPE TYPE; /* . . . */
};

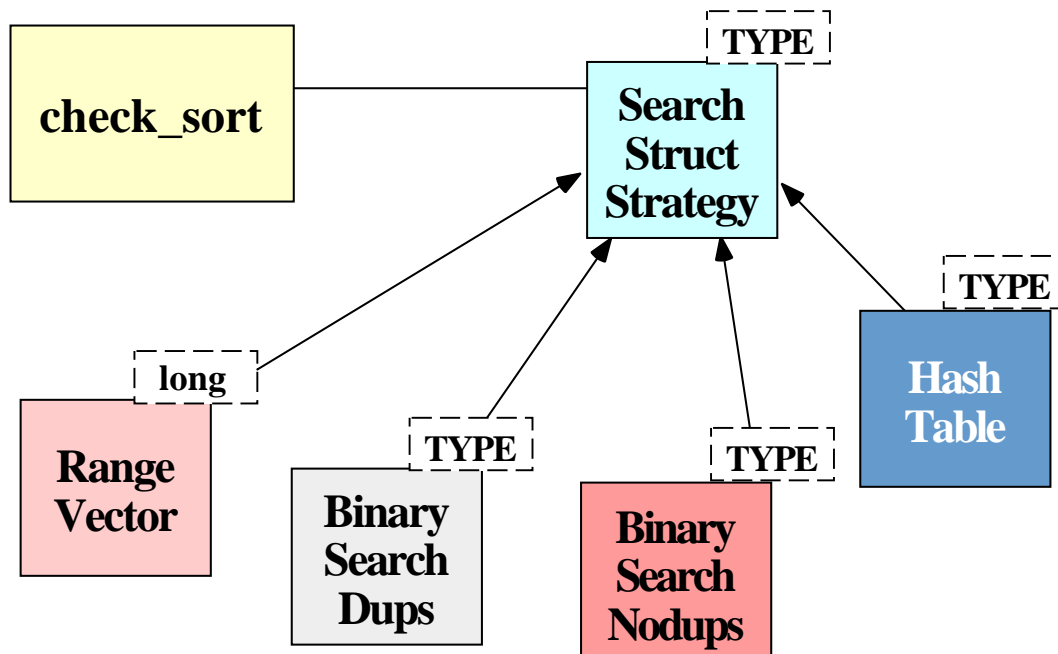
template <typename T>
class Hash_Table :
public Search_Strategy<T>
{
    typedef typename ARRAY::TYPE TYPE; /* . . . */
};
```

---

## Design Patterns in Sort Verifier

- Factory Method
  - *Define an interface for creating an object, but let subclasses decide which class to instantiate*
    - \* Factory Method lets a class defer instantiation to subclasses
- In addition, the *Facade*, *Iterator*, *Singleton*, & *Strategy* patterns are used

## Using the Strategy Pattern



- This pattern extends the strategies for checking if an array is sorted without modifying the `check_sort` algorithm

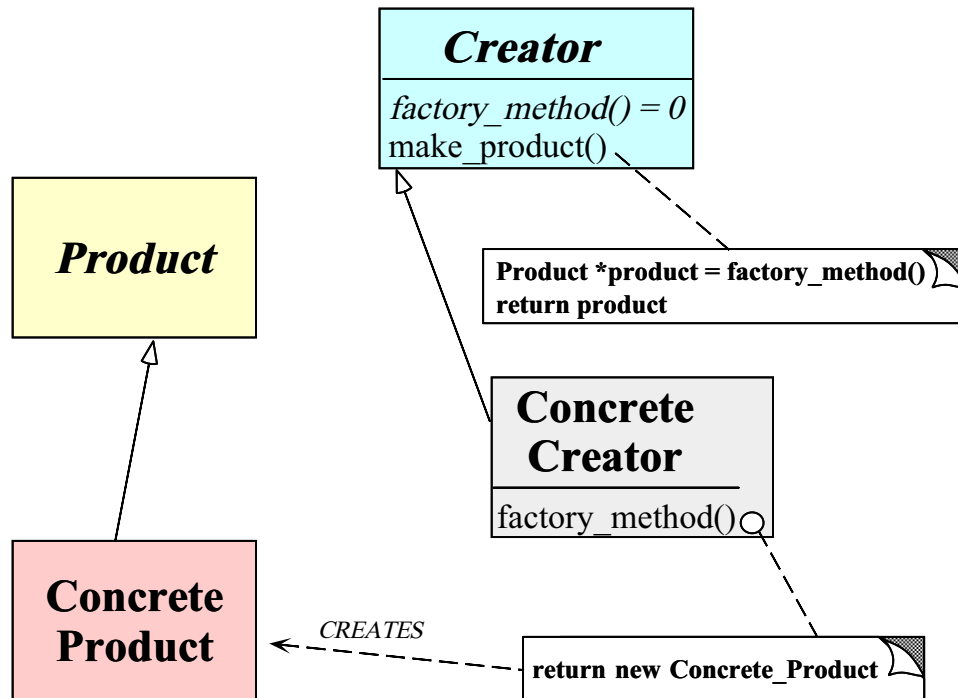


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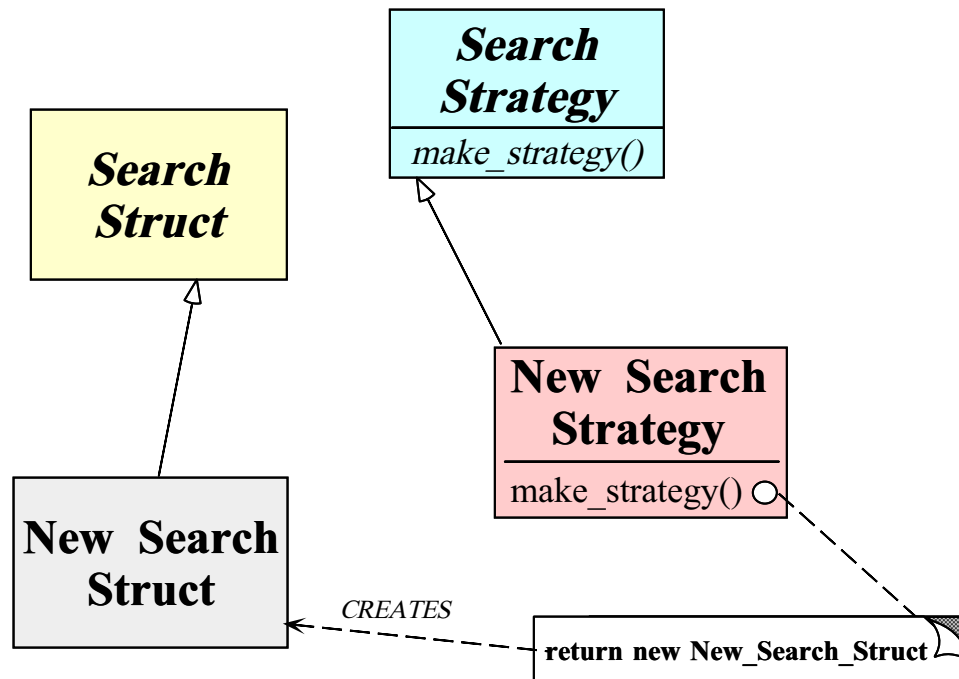
## The Factory Method Pattern

- *Intent*
  - Define an interface for creating an object, but let subclasses decide which class to instantiate
    - \* Factory Method lets a class defer instantiation to subclasses
- This pattern resolves the following force:
  1. *How to extend the initialization strategy in the sort verifier transparently*

## Structure of the Factory Method Pattern



## Using the Factory Method Pattern



---

## Implementing the check\_sort Function

- e.g., C++ code for the sort verification strategy

```
template <typename ARRAY> int
check_sort (const ARRAY &orig,
            const ARRAY &p_sort) {
    if (orig.size () != p_sort.size ())
        return -1;

    auto_ptr < Search_Strategy<typename ARRAY::TYPE> > ss =
        Search_Struct<ARRAY>::instance ()->make_strategy
        (p_sort);
```

## Implementing the check\_sort Function (cont'd)

```
for (int i = 0; i < p_sort.size (); ++i)
    if (ss->insert (p_sort[i]) == false)
        return -1;

for (int i = 0; i < orig.size (); ++i)
    if (ss->remove (orig[i]) == false)
        return -1;

return 0;
// auto_ptr's destructor deletes the memory . . .
}
```

## Initializing the Search Structure

- Factory Method

```
template <typename ARRAY>
Search_Strategy<typename ARRAY::TYPE> *
Search_Struct<ARRAY>::make_strategy
    (const ARRAY &potential_sort) {
    int duplicates = 0;

    for (size_t i = 1; i < potential_sort.size (); ++i)
        if (potential_sort[i] < potential_sort[i - 1])
            return 0;
        else if (potential_sort[i] == potential_sort[i - 1])
            ++duplicates;
```

## Initializing the Search Structure (cont'd)

```
if (typeid (potential_sort[0]) == typeid (long)
    && range <= size)
    return new Range_Vector (potential_sort[0],
                             potential_sort[size - 1])
else if (duplicates == 0)
    return new Binary_Search_Nodups<ARRAY>
           (potential_sort);
else if (size % 2)
    return new Binary_Search_Dups<ARRAY>
           (potential_sort, duplicates)
else return new Hash_Table<typename ARRAY::TYPE>
           (size, &hash_function);
}
```

## Specializing the Search Structure for Range Vectors

```
template <Array<long> > Search_Strategy<long> *  
Search_Struct<Array<long> >::make_strategy  
    (const Array<long> &potential_sort)  
{  
    int duplicates = 0;  
  
    for (size_t i = 1; i < size; ++i)  
        if (potential_sort[i] < potential_sort[i - 1])  
            return 0;  
        else if (potential_sort[i] == potential_sort[i - 1])  
            ++duplicates;  
  
    long range = potential_sort[size - 1] -  
                potential_sort[0];
```



## Specializing the Search Structure for Range Vectors

```
if (range <= size)
    return new Range_Vector (potential_sort[0],
                             potential_sort[size - 1])
else if (duplicates == 0)
    return new Binary_Search_Nodups<long>
           (potential_sort);
else if (size % 2)
    return new Binary_Search_Dups<long>
           (potential_sort, duplicates)
else return new Hash_Table<long>
           (size, &hash_function);
}
```

---

## Summary of Sort Verifier Case Study

- The sort verifier illustrates how to use OO techniques to structure a modular, extensible, & efficient solution
  - The main processing algorithm is simplified
  - The complexity is pushed into the strategy objects & the strategy selection factory
  - Adding new solutions does not affect existing code
  - The appropriate ADT search structure is selected at run-time based on the Strategy pattern