

The neglected art of Fixed Point arithmetic

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Contents

- Motivation
- Introduction
- Typically needed functions
- Caveats and tricks
- Tips for making a fixed point library

Motivation

- *Man sent himself to moon, and space probes even beyond that.* Do you think the hardware used to accomplish those feats had fancy FPU to do all the calculations?
- They used RCA 1802.
 - Processing power equals roughly 6502 or 6510, used in Apple II and Commodore 64.

Motivation

- But it's a lot of work.
 - 30% of the Apollo software development effort was spent on scaling. [KrL64]
- So they eventually switched to floating point when hardware got better.

Motivation

- So why am I talking about this?
 - Well, at least it's COOL, in retro-way:
This is how demo & game coders did their 3D stuff 15 years ago and made some pretty cool stuff even with the minuscule CPU power.
- But does that matter anymore - except if you are going to take part in the old school demo competition with some retro stuff?

Motivation

- There's still plenty of platforms where using only fixed point (integer) calculations is still very relevant.
 - Mobile devices (Typical: ARM CPU, no FPU)
 - Almost all mobile phones (J2ME or native code)
 - Handheld consoles (Gameboy, Nintendo DS)
 - DSP Programming
 - There's both fixed & floating point DSPs

Motivation

- ...continued...
 - OpenGL ES is the standard for embedded 3D.
 - Profiles for both fixed point and floating point, but often only Common-Lite profile is provided (no floating point).
 - Fixed point is often still a bit faster on desktop than floating point.
 - Stable calculations across platforms
 - Floating point calculations are prone to slight differences based on compiler, CPU and other dependencies.

Introduction

- Basics
- Notation
- Range and precision
- Conversion
- Basic operations: + - * /

Introduction: Basics

- What are the fixed point numbers in “layman's” terms?
 - Scale all real numbers by a constant factor, such as 65536, round to nearest integer and store the numbers as integers.
 - This allows you to represent an evenly distributed subset of real numbers roughly from -32768 to 32767 (with 32-bit signed integers and factor of 65536).

Introduction: Basics

- More exactly, you are dividing your range of values to two parts - the integer part and fractional part.

2^7	2^6	2^5	2^4	2^3	2^2	2^1	2^0
128	64	32	16	8	4	2	1

8-bit example:

2^3	2^2	2^1	2^0	.	2^{-1}	2^{-2}	2^{-3}	2^{-4}
8	4	2	1	.	1/2	1/4	1/8	1/16

That's the "*fixed point!*"

Introduction: Notation

- Notations:
 - $M.N$, e.g. 16.16
 - QN (Q factor), e.g. Q16
- M is number of integer bits and N is number of fractional bits.

Introduction:

Range and precision

- Range: defined by the integer (upper) part.
 - 16.16 (signed): range is [-32768, 32767]
- Precision: smallest difference between two successive numbers is $1/2^N$.
 - 16.16: $1/65536$ (~0.000015258789)

8-bit example:

2^3	2^2	2^1	2^0	.	2^{-1}	2^{-2}	2^{-3}	2^{-4}
8	4	2	1	.	1/2	1/4	1/8	1/16

4.4

Range: [-8, 7]
(if signed)

Precision: 1/16
(0.0625)

Introduction: Conversion

- Conversion from real to fixed point number
 - Multiply by 2^N and round to nearest integer
 - $(\text{int}) (\mathbf{R} * (1 \ll \mathbf{N}) + (\mathbf{R} >= 0 ? 0.5 : -0.5))$
- Conversion from fixed point to real number
 - Cast to real and divide by 2^N
 - $(\text{float}) \mathbf{F} / (1 \ll \mathbf{N})$
- Conversion from/to integers (lossless)
 - Shift N bits up or down (scaling by 2^N)
 - $F = I \ll N, \quad I = F \gg N$

Introduction:

Basic operations

- Addition (+) and subtraction (-)
 - Same as adding and subtracting integers
- Multiplication ($a * b$)
 - Multiply as integers and divide result by 2^N .
 - $((a * b) \ggg N)$
 - That overflows very easily, as both a and b are fixed point numbers!
 - If both a and b are 2.0 (131072) as 16.16 fixed point
 $(a * b) == 17179869184$ - 32 bits isn't enough!

Introduction: Basic operations

- For multiplication, the intermediate result from $(a * b)$ is in $2M:2N$ ($Q2N$) format
 - Store intermediate value in double sized integer format. That is, for 32-bit 16.16 fixed point numbers, you need a 64-bit integer to store the 32.32 ($Q32$) intermediate result.

- $(int) ((INT64) a * (INT64) b) >> N)$

INT64
MSVC: `__int64`
GCC: `signed long long`
Java: `long`

Introduction: Basic operations

- Division (a / b)
 - Multiply a by 2^N and divide by b (as integers).
 - ~~$((a \ll N) / b)$~~
 - Again, intermediate result is prone to overflowing, so the correct way for 16.16 is:
 - $(((INT64) a \ll N) / b)$
- See references for more detailed introductory texts to fixed points. [VVB04, Str04, WikF]

Typically Needed Functions

- Sine and cosine: $\sin(x)$, $\cos(x)$
- Arcus tangent: $\text{atan2}(y, x)$
- Square root: $\text{sqrt}(x)$
- Try CORDIC

Typically Needed Functions: Sine and cosine

- Typical approach is to use a look-up table.
 - Requires memory proportional to desired accuracy
 - Requires some storage space to load table from or time for pre-calculating table on startup
 - Can interpolate between sampled values to gain some more accuracy
- Note that it's enough to calculate $\pi/4$ entries to table, rest of the samples can be mirrored and transformed from those.

Typically Needed Functions: Sine and cosine

- It's possible to find or construct less accurate approximations for functions if you need smaller code, memory usage or more speed.
 - DSP coders have some quite nice tricks. [Ben06]
- See also [Str04] for code example of how to calculate sin, cos and tan algorithmically using only a small arctan table.

Typically Needed Functions: Square root

- Several fairly good iterative algorithms exist, so I don't recommend using a look-up table.
- Can be as simple as trying out to multiply integers by themselves until you find out the closest one
 - Or binary search version of the above
- Ken Turkowski's implementation is probably the most often used one. [Tur94]
 - For your convenience, code on the next slide.

Typically Needed Functions: Square root

```
/* The definitions below yield 2 integer bits, 30 fractional bits */
#define FRACBITS 30      /* Must be even! */
#define ITERS      (15 + (FRACBITS >> 1))
typedef long TFract;

TFract
FFracSqrt(TFract x)
{
    register unsigned long root, remHi, remLo, testDiv, count;

    root = 0;           /* Clear root */
    remHi = 0;          /* Clear high part of partial remainder */
    remLo = x;          /* Get argument into low part of partial remainder */
    count = ITERS;      /* Load loop counter */

    do {
        remHi = (remHi << 2) | (remLo >> 30); remLo <<= 2; /* get 2 bits of arg */
        root <<= 1; /* Get ready for the next bit in the root */
        testDiv = (root << 1) + 1; /* Test radical */
        if (remHi >= testDiv) {
            remHi -= testDiv;
            root += 1;
        }
    } while (count-- != 0);

    return(root);
}
```

[Tur94]

Typically Needed Functions: Arcus tangent

- You can try some look-up table tricks, again.
- If fast and rough approximation is enough, implementation can be very simple. [Cap91]
- For accurate results, try using CORDIC (covered next).
- For my favorite approximation (for the time being), check Jim Shima's DSP Trick: Fixed-Point Atan2 With Self Normalization. [Shi99]

Typically Needed Functions: Try CORDIC

- “COordinate Rotation Digital Computer”, an algorithm to calculate hyperbolic and trigonometric functions, from 1959. [WikC]
 - Only small look-up tables, bitshifts and additions.
- Use it run-time or to pre-calculate look-up tables. (sin, cos, atan, ...)
- Accurate results
- Not the fastest solution

Caveats And Tricks

- Back to range and precision
- Watch out for division by zero
- Exact results
- Dealing with problems

Caveats And Tricks:

Back to range and precision

- When storing result of $a*b$ to normal sized fixed point (integer) value
 - Possible range & precision for the original values is much more limited than the normal to prevent overflow & underflow.
 - For storing $a*a$:
 - $\text{abs}(a) \leq \sim 181$ -- $181*181 = 32761$, barely fits in signed 16.16 fixed point number.
 - $\text{abs}(a) \geq \sim 0.004$ -- $0.004*0.004 = 0.000016$, truncated down to $1/65536$.

Caveats And Tricks:

Back to range and precision

- Similarly, make sure that a/b will stay in range
 - When $|b| > 1.0$
 - Check ranges so that result doesn't end up being 0.
 - When $|b| < 1.0$
 - $b > 1 / (2^{M-1} / a)$
 - If max value for a is 32, b must be at least 0.000991821 (65/65536) so that a/b fits in 16.16 fixed point number: $32 / 0.000991821 \approx 32263$.
 - If b would be one less (64/65536), then a/b will be 32768, not fitting in $[-32768, 32767]$ 16.16 fixed point value range.

Caveats And Tricks: Watch out for division by zero

- Floating points have “Infinity Arithmetic”
 - Even result of division by zero is defined, so you simply get Inf as a result
 - Easier to go unnoticed by mistake
- Fixed point (integer) division by zero leads to interrupt or an exception is thrown
 - Typically programs just crash at this

Caveats And Tricks: Exact results

- Possible in some cases: modify division involving formulas to keep numerator and denominator separate, and try to find out final (exact) result by examining those, without doing the division. See [Eri05] for example.
- Generally speaking, it's rare and hard to take advantage of this.

Caveats And Tricks: Dealing with problems

- When troubled by overflows, underflows or accuracy problems
 - Try keeping the intermediate result(s) in the bigger (64 bit) format and work out the final result directly from there.
 - Use asserts and do other verification checks rigorously, especially in debug builds.
 - Compare to results of same calculations done in floating points.

Tips For Making A Fixed Point Library

- There's built-in support... if you code in Ada.
- C/C++ alternatives:
 - Code it all in-line, using normal integers
 - Use helper macros (conversions, operations)
 - Create a real number class with overloaded operators
 - Allows to switch easily between floats and fixed points

Tips For Making A Fixed Point Library

- Create debug version of the real number class
 - Perform both fixed point and floating point calculations in parallel
 - Detect overflow & underflow conditions
 - Detect drifting
 - Error/warning asserts and checks can be made run-time togglable
- If you work on J2ME, it's best to inline all calculations yourself for performance.

Other Tidbits

- Nobody noticed that I changed the underlying physics engine from floating point to fixed point in latest version of *Pogo Sticker*.
- You can do fixed point (integer) `abs()` without branches. [And05, War02]
 - For 32-bit ints:
 - `result = (v ^ (v >> 31)) - (v >> 31)`
 - Ridiculously that's patented. But that's not the only way, check the references.

Other Tidbits

- 32-bit signed 0x80000000 (highest bit) is special
 - `int x; if (x < 0) x = -x;`
Doesn't work as expected if `x==0x80000000!`
X will still be 0x80000000 (-2147483648).
 - For the above example, solution is to cast result to unsigned int as you know it will not be negative.

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Thank You!

- Questions & Answers
 - If there's time
- Slides will be available from my home page:
 - <http://jet.ro>
- Get some games:
 - <http://www.skinflake.com>